

Atelier B

Reusable Components

Reference Manual

version 3.6



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Reusable Components—Reference Manual
version 3.6

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Introduction

The reusable components supplied with Atelier B are basic machines and library machines. Basic machines are the modelisation in B of modules manually coded in C, C++ or ADA. These modules are used to encapsulate the operating system functions that must be used; they must usually be performed in taking into account the specificities of the hardware that the security software will run on. This is why there are few basic machines delivered with "Atelier B".

Library machines are abstract machines written in B language. They generally model a type of mathematical object (sequence, function, etc.) and offer the operations that allow the handling of these objects.

Unlike basic machines, library machines are properly performed using the B method, i.e., using refining and implementation in B along with complete proof of the set. This proof may in principle be executed at any time in order to check its validity (warning: proving methods may depend on the demonstrator version used). Therefore, unlike basic machines, library machines may be numerous and complex while remaining secure as they are proven.

To use basic machines, simply reference them in the appropriate B project, by INCLUDES, IMPORTS or any derived actions. When the final project is translated into a traditional programming language, the translation of the library machine implementations used must be redone if this was not already done at Atelier B installation.

Library machines are implemented on basic machines. As they are performed until the implementation in B language, they provide complete examples of use in the B method. They especially contain examples of proven WHILE loops. For practical advice on proving WHILE loops, refer to the "B Language User Manual".

The user may directly use library machines just like he uses basic machines. Sometimes the implementation of a library machine may use the services of a machine that it does not create an instance for (use by SEES) to avoid duplications. In this case the user will have to create the instance in question (using IMPORTS) by following the indications in the "IMPORTS REQUIRED" section of the description for each library machine.

When the final project C, C++ or ADA compilation is performed, the library compilation is automatically performed if necessary. Performing link editing will then enable incorporation into the final executable program only those object files that correspond to the library machines actually used. All this is performed in the *Makefile* produced by Atelier B. To integrate a software component produced by Atelier B into a traditional product, use this *Makefile* as a basis or refer to the "ADA Translator User Manual".

Warning:

This warning regards the use of reusable components with the Ada, C and C++ translators

supplied with Atelier B. These translators are experimentals. Their goal is to show that it is possible to translate some B0 implementations into classical programming languages. Therefore, their use is not guaranteed. Especially the reusable components use may induce errors when compiling the code produced by the translators. The reusable components must be considered as examples. Each user can develop his own library machines according to his needs.

2 Index of Basic Machines

BASIC_ARRAY_VAR: implanting a one dimension table

VAL_ARRAY read a table element

STR_ARRAY write a table element

BASIC_ARRAY_RGE: implementing a two dimensional table

VAL_ARR_RGE read a table element

STR_ARR_RGE write a table element

COP_ARR_RGE copy a table line to another

CMP_ARR_RGE compare two table lines

BASIC_IO: vt 100 style input/output

INTERVAL_READ entry by the operator of a number in mm..nn.

INT_WRITE print a number.

BOOL_READ entry by a TRUE or FALSE boolean operator

BOOL_WRITE print the TRUE or FALSE condition.

CHAR_READ entry by a character's operator.

CHAR_WRITE print a character.

STRING_WRITE print a message.

3 Index of Library Machines

L_ARITHMETIC1: extended integer operations: MIN, MAX, INC, DEC, EXP, SQRT, LOG

VAL_ARR_RGE read a table element

STR_ARR_RGE write a table element

COP_ARR_RGE copy a table line to another

CMP_ARR_RGE compare two table lines

BASIC_IO: vt 100 style input/output

MIN minimum of two numbers.

MAX maximum of two numbers.

INC increment a number.

DEC decrement a number.

EXP exponential.

SQRT integer square root by default.

LOG_BY_DEFAULT logarithm by default.

LOG_BY_EXCESS logarithm by excess.

L_ARRAY1: one dimensional table with initialization loop

VAL_ARRAY value of an element (promoted operation)

STR_ARRAY write an element (promoted operation)

SET_ARRAY write the same value in a portion of the table

L_ARRAY3: table with non-ordered values, maximum operations

VAL_ARRAY value of an element (promoted operation).

STR_ARRAY write an element (promoted operation).

SET_ARRAY write a same value in a table portion (promoted operation).

SWAP_ARRAY exchange two elements (promoted operation).

RIGHT_SHIFT_ARRAY shift a portion to the large index (promoted operation).

LEFT_SHIFT_ARRAY shift a portion to the small index (promoted operation).
 SEARCH_MAX_EQL_ARRAY search for a value in a portion of the table (promoted operation).
 SEARCH_MIN_EQL_ARRAY search for a value in a portion of the table (promoted operation).
 REVERSE_ARRAY invert the order of the elements in a portion of the table.

L_ARRAY5: table with ordered values, sort operation

VAL_ARRAY value of an element (promoted operation).
 STR_ARRAY write an element (promoted operation).
 SET_ARRAY write the same value in a portion of the table (promoted operation).
 SWAP_ARRAY exchange two elements (promoted operation).
 RIGHT_SHIFT_ARRAY shift a portion to the large index (promoted operation).
 LEFT_SHIFT_ARRAY shift a portion to the small index (promoted operation).
 SEARCH_MAX_EQL_ARRAY search for a value in a portion of the table (promoted operation).
 SEARCH_MIN_EQL_ARRAY search for a value in a portion of the table (promoted operation).
 REVERSE_ARRAY invert the order of elements in a portion of the table (promoted operation).
 SEARCH_MIN_GEQ_ARRAY search for the first element that exceeds a value (promoted operation).
 ASCENDING_SORT_ARRAY sort of a table portion.

L_PFNC: partial function

VAL_PFNC value of the function for an element in its domain
 STR_PFNC overloads the partial function with a couple
 XST_PFNC tests if an index is in the partial function domain
 RMV_PFNC removes a couple from the partial function
 SET_PFNC overloads a part of the function with a constant
 SWAP_PFNC exchanges the images for two domain indexes
 RIGHT_SHIFT_PFNC right shift of a domain part
 LEFT_SHIFT_PFNC left shift of a domain part
 SEARCH_MAX_EQL_PFNC searches for a value in the partial function
 SEARCH_MIN_EQL_PFNC searches for a value in the partial function
 REVERSE_PFNC reverses the order of elements for a portion of the domain

ASCENDING_SORT_PFNC sorts in a portion of the domain

L_SEQUENCE: building a sequence

LEN_SEQ returns the current size of the sequence.
IS_FULL_SEQ is used to determine if the sequence is full (size = LS_maxsize).
IS_INDEX_SEQ is used to determine whether ii is a valid index.
VAL_SEQ value of an element in the sequence.
FIRST_SEQ returns the first element in the sequence.
LAST_SEQ returns the last element in the sequence.
PUSH_SEQ add vv to the end of the sequence.
POP_SEQ removes the last element from the sequence (its value is lost).
STR_SEQ changes the value of an element in the sequence.
RMV_SEQ removes an element from the middle of the sequence.
INS_AFT_SEQ inserts vv right after index ii.
CLR_SEQ clears the sequence.
TAIL_SEQ removes the first element from the sequence.
KEEP_SEQ only keeps the first elements in the sequence.
CUT_SEQ cuts the nn first elements from the sequence.
PART_SEQ only retains part ii..jj in the sequence.
REV_SEQ reverses the order of elements in the sequence.
FIND_FIRST_SEQ finds vv in the sequence, from the start.
FIND_LAST_SEQ finds vv in the sequence, from the end.

L_SET: creating a set

CARD_SET returns the cardinal for the set.
IS_FULL_SET identifies if the set is full (card = LSET_maxsize).
IS_INDEX_SET identifies if a number is a valid index.
VAL_SET value of a element in the set.
FIND_SET finds an element in the set.
RMV_SET removes an element from the set.
INS_SET inserts an element in the set.
CLR_SET clears all elements from the set.

L_ARRAY_1_RANGE: array of tables of the same size with numerical indexes

VAL_ARR_RGE value of an element (promoted operation).

STR_ARR_RGE	write an element (promoted operation).
COP_ARR_RGE	copy a table to another (promoted operation).
CMP_ARR_RGE	compare two tables (promoted operation).
DUP_ARR_RGE	duplicate the same table into a series of tables.
SET_ARR_RGE	copy the same value to an index set in one of the tables.
PCOP_ARR_RGE	copy part of one of the tables to a different table to a given position.
PCMP_ARR_RGE	find the first element that is different from two parts of two tables. A Boolean element indicates if this element was found and, in this case, the index of this element is returned.

L_ARRAY_3_RANGE: range of tables of the same size, with numerical indexes, and values that are not ordered, maximum operations

VAL_ARR_RGE	value of an element (promoted operation).
STR_ARR_RGE	write an element (promoted operation).
COP_ARR_RGE	copy a table to another (promoted operation).
CMP_ARR_RGE	compare two tables (promoted operation).
DUP_ARR_RGE	duplicate the same table to an array of tables (promoted operation).
SET_ARR_RGE	copy the same value to a range in one of the tables (promoted operation).
PCOP_ARR_RGE	copy part of one of the tables to a different table, in a given position (promoted operation).
PCMP_ARR_RGE	find the first different element from two parts in two tables. A Boolean element indicates whether this element was found and, in this case, the index of this element is returned (promoted operation).
SWAP_RGE	swap two elements in a table.
RIGHT_SHIFT_RGE	shift a table range to the large index.
LEFT_SHIFT_RGE	shift a table range to the small index.
SEARCH_MAX_EQL_RGE	find the last element that equals a value in a table range.
SEARCH_MIN_EQL_RGE	find the first element that equals a value in a table range.
REVERSE_RGE	reverse the order of the elements of a table part.

L_ARRAY_5_RANGE: array of tables of the same size, with numerical indexes, with ordered values, sort operations

VAL_ARR_RGE	value of an element (promoted operation).
STR_ARR_RGE	write an element (promoted operation).

COP_ARR_RGE	copy a table to another (promoted operation).
CMP_ARR_RGE	compare two tables (promoted operation).
DUP_ARR_RGE	duplicate the same table in a range of tables (promoted operation).
SET_ARR_RGE	copy the same value to an index range in one of the arrays (promoted operation).
PCOP_ARR_RGE	copy a range from one of the tables to a different table, at a given position (promoted operation).
PCMP_ARR_RGE	find the first different element in two ranges in two tables. A Boolean element indicates that this element was found and, in this case, the index of this element is returned (promoted operation).
SWAP_RGE	swap two elements in a table (promoted operation).
RIGHT_SHIFT_RGE	shift a table range to the large index (promoted operation).
LEFT_SHIFT_RGE	shift a table range to the small index (promoted operation).
SEARCH_MAX_EQL_RGE	search for the last element that equals a value in a table range (promoted operation).
SEARCH_MIN_EQL_RGE	search for the first element that equals a value in a table range (promoted operation).
REVERSE_RGE	reverses the order of the elements of a part of a table (promoted operation).
SEARCH_MIN_GEQ_RGE	search for the first element that exceeds a value in a table range.
ASCENDING_SORT_RGE	sort a table range into ascending order.

L_SEQUENCE_RANGE: sequence range

LEN_SEQ_RGE	determines the length of a sequence.
IS_FULL_SEQ_RGE	determines whether a sequence is full.
IS_INDEX_SEQ_RGE	determines whether an integer is in a sequence range.
VAL_SEQ_RGE	gives the value of a sequence for a valid index.
FIRST_SEQ_RGE	gives the first element in a sequence.
LAST_SEQ_RGE	gives the last element in a sequence.
PUSH_SEQ_RGE	adds an element to a sequence.
POP_SEQ_RGE	removes the last element from a sequence.
STR_SEQ_RGE	changes the value of a sequence element.
RMV_SEQ_RGE	removes an element from a sequence, with a size that decreases by 1.
INS_SEQ_RGE	adds an element to a sequence, with a size that increases by 1.
CLR_SEQ_RGE	clears a sequence.
TAIL_SEQ_RGE	removes the first element from a sequence.

KEEP_SEQ_RGE	only keeps in a sequence the N first elements.
CUT_SEQ_RGE	cuts the N first elements from a sequence.
PART_SEQ_RGE	only keeps in a sequence the indexes in a range between two limits.
REV_SEQ_RGE	reverses the order of the elements in a sequence.
FIND_FIRST_SEQ_RGE	finds a value in a sequence, returns a Boolean element indicating that it was found and if yes returns the smallest corresponding index.
FIND_LAST_SEQ_RGE	finds a value in a sequence, returns a Boolean element indicating that it was found and if yes returns the largest corresponding index.
COP_SEQ_RGE	copies from one sequence to another.
CMP_SEQ_RGE	comparison of two sequences.
PCOP_SEQ_RGE	partially copies one of the sequences to another.
PCMP_SEQ_RGE	partial comparison of two sequences.

L_ARRAY_COLLECTION: collection of arrays of the same size

CRE_ARR_COL	returns a Boolean element indicating that there is still an array free in the collection and gives the index of this free array.
DEL_ARR_COL	releases the identified array.
VAL_ARR_COL	reads an element from one of the valid arrays.
STR_ARR_COL	writes an element from one of the valid arrays.
COP_ARR_COL	copies one of the arrays to another.
CMP_ARR_COL	compares two tables.

L_ARRAY1_COLLECTION: collection of arrays of the same size with numerical index

CRE_ARR_COL	returns a Boolean element indicating that there is an array free in the collection and the index of this free array (promoted operation).
DEL_ARR_COL	releases the listed array (promoted operation).
VAL_ARR_COL	read a element from on of the valid arrays (promoted operation).
STR_ARR_COL	write a element from one of the valid arrays (promoted operation).
COP_ARR_COL	copies from one of the arrays to another (promoted operation).
CMP_ARR_COL	compares two tables (promoted operation).
SET_ARR_COL	copies the same value to an index range in one of the arrays.
PCOP_ARR_COL	copies part of one of the arrays to another, to a given position.
PCMP_ARR_COL	find the first different element between the two parts of the two different arrays. A Boolean element indicates if this element was found and in this case, the index of this element is returned.

L_RELATION : complete binary relations

op_reset	The relation becomes the empty relation.
op_isFullRelation	Returns TRUE only if the cardinal of the relation equals <code>max_nb- _2tuple</code> .
op_add	Adds a couple to the relation.
op_remove	Removes a couple to the relation.
op_cardinal	Returns the relation cardinal ¹
op_belongsTo	Checks if a couple is present in the relation.

¹i.e. the number of couple present in the relation.

4 Description of Basic Machines

The basic machines supplied with Atelier B allow either the creation of dynamic arrays that cannot be obtained using B0, or producing models using vt100 style inputs/outputs. “dynamics arrays” are arrays which size depends on the machine parameters. Such arrays cannot be realised directly in B0, the safety design of the ADA, C and C++ translators do not allow to treat this case. For example, the following construction is not allowed:

```

IMPLEMENTATION
  mm(xx)
...
CONCRETE_VARIABLES
  mytab
INVARIANT
  mytab ∈ (0..xx) → NAT
...
END

```

Such an array would have to be realised using BASIC_ARRAY_VAR.

The atelier actual version is composed of three basic machines:

BASIC_ARRAY_VAR Arrays with dimension 1.

BASIC_ARRAY_RGE Arrays with dimension 2.

BASIC_IO Usual inputs/outputs management.

This chapter presents this three machines.

The basic machine BASIC_IO is intended to the model designing. It mustn't be considered as safe.

WARNING: The manual implementations of the basic machines BASIC_ARRAY_VAR and BASIC_ARRAY_RGE destined for the translators supplied with Atelier B are provided as demonstration. They are not safe, and are not appropriated in all the B use contexts.

4.1 BASIC_ARRAY_VAR: Implanting a one dimensional table

OPERATIONS

VAL_ARRAY read a table element
STR_ARRAY write a table element

EXAMPLE

Example of use with listed sets:

<pre> MACHINE array SETS FONTS = {Times,Serif,Courier}; FTYPE = {fixed,unfixed} VARIABLES fixedsz INVARIANT fixedsz ∈ FONTS → FTYPE INITIALISATION fixedsz:={Times ↦ unfixed, Serif ↦ fixed, Courier ↦ fixed} END </pre>	<pre> IMPLEMENTATION array_1 REFINES array IMPORTS BASIC_ARRAY_VAR(FONTS,FTYPE) INVARIANT arr_vrb = fixedsz INITIALISATION STR_ARRAY(Times,unfixed); STR_ARRAY(Serif,fixed); STR_ARRAY(Courier,fixed) END </pre>
--	--

arr_vrb is the name of the table encapsulated by BASIC_ARRAY_VAR

DESCRIPTION

BASIC_ARRAY_VAR modelizes one dimensional arrays. Such arrays cannot be created directly in B0 if their size depend on the machine parameters (“dynamic arrays”). The current design of ADA or C translators does not allow handling this case. The following construction is therefore illegal:

```

IMPLEMENTATION
  mm(xx)
...
VARIABLES
  mytab
INVARIANT
  mytab ∈ (0..xx) → NAT
...
END

```

This kind of table should be generated using BASIC_ARRAY_VAR.

MACHINE PARAMETERS

BASIC_ARRAY_VAR (BAV_INDEX,BAV_VALUE): BAV_INDEX is the set of values used to index the table, BAV_VALUE is the set of possible values for table elements.

The B language rule relating to the possible values of the BAV_VALUE parameter ensure that: if a computer variable can contain elements of MININT..MAXINT, then it can contain those of BAV_VALUE. For example, B rules forbid assigning BAV_VALUE the value of MAXINT+1,MAXINT+2

VAL_ARRAY

syntax $vv \leftarrow \text{VAL_ARRAY}(ii)$

preconditions ii must be a BAV_INDEX

outputs vv is a BAV_VALUE, the value of the array at position ii .

STR_ARRAY

syntax $\text{STR_ARRAY}(ii, vv)$

preconditions ii must be a BAV_INDEX and vv must be a BAV_VALUE

The value vv is stored in the array at ii index.

C++ LANGUAGE

In C++, the array is realised by an integer array. The accesses to this array are done using method that refuse the index used between 0 and the array size, guaranting an optimal memory use.

The array is dynamically reserved when launching the program. If the size indicated by the formal parameters is too big, the program stops with the following message:

Virtual memory `exceede` in `‘‘new’’`

C LANGUAGE

The realisation in C is based on the same principles as in C++. The stop message on initial reservation failure is:

Fatal error: Malloc of X bytes failed
Execution of current application is aborted

ADA LANGUAGE

The use of generic packaging guarantees an optimal memory occupation. No restrictions are made on the instantiation parameters. On initial reservation failure, an exception stops the program.

PROGRAMMING

Example of use with literal sets:

<pre>MACHINE narr VARIABLES myvar INVARIANT myvar ∈ 0..2 → 0..1 INITIALISATION myvar:= {0 ↦ 0, 1 ↦ 1, 2 ↦ 1} END</pre>	<pre>IMPLEMENTATION narr_1 REFINES narr IMPORTS BASIC_ARRAY_VAR(0..2,0..1) INVARIANT arr_vrb = myvar INITIALISATION STR_ARRAY(0,0); STR_ARRAY(1,1); STR_ARRAY(2,1) END</pre>
---	--

Another example. Only the implementation is presented. The write of a machine refined by this implementation is an exercice for the reader:

<pre>IMPLEMENTATION parr_1 REFINES parr IMPORTS BASIC_ARRAY_VAR(FONTS,FTYPE) VALUES FONTS = 5..7; FTYPE = 3..4 INVARIANT arr_vrb = fixedsz INITIALISATION STR_ARRAY(5,3); STR_ARRAY(6,4); STR_ARRAY(7,5) END</pre>

NOTE: The possible values of the BASIC_ARRAY_VAR parameters are given by the B language rules, (refer to section 12.2 page 574 of the BBOOK)

4.2 BASIC_ARRAY_RGE: Implementing a Two Dimensional Array

OPERATIONS

- VAL_ARR_RGE read an array element
- STR_ARR_RGE write an array element
- COP_ARR_RGE copy an array line to another
- CMP_ARR_RGE compare two array lines

EXAMPLE

Example of use, two lines and three columns array:

<p>MACHINE bitab</p> <p>SETS LGNS = {ll1,ll2}</p> <p>VARIABLES mytab</p> <p>INVARIANT mytab ∈ LGNS → (1..3 → 0..255)</p> <p>INITIALISATION mytab:={ll1 ↦ {1 ↦ 7,2 ↦ 8,3 ↦ 9}, ll2 ↦ {1 ↦ 0,2 ↦ 1,3 ↦ 2}}</p> <p>END</p>	<p>IMPLEMENTATION bitab_1</p> <p>REFINES bitab</p> <p>IMPORTS BASIC_ARRAY_RGE(1..3,0..255,LGNS)</p> <p>INVARIANT arr_rge = mytab</p> <p>INITIALISATION STR_ARR_RGE(ll1,1,7); STR_ARR_RGE(ll1,2,8); STR_ARR_RGE(ll1,3,9); STR_ARR_RGE(ll2,1,0); STR_ARR_RGE(ll2,2,1); STR_ARR_RGE(ll2,3,2)</p> <p>END</p>
---	--

The variable arr_rge is the name of the encapsulated array par BASIC_ARRAY_RGE

DESCRIPTION

BASIC_ARRAY_RGE models two dimensional arrays. Such arrays cannot be created directly in B0 if their size depends on the machine parameters (“dynamic array”). The safe design of the ADA, C++ or C translators do not allow to treat this case. The following construction is forbidden:

<p>IMPLEMENTATION mm(xx) ...</p> <p>CONCRETE_VARIABLES mytab</p> <p>INVARIANT mytab ∈ (0..10) → (0..xx) × (0..xx)</p> <p>...</p> <p>END</p>

Such an array must be implemented using BASIC_ARRAY_RGE.

MACHINE PARAMETERS

BASIC_ARRAY_RGE(BAR_INDEX,BAR_VALUE,BAR_RANGE):

BAR_INDEX represents the column indexes.

BAR_VALUE is the set of the possible values for the array elements,

BAR_RANGE represents the line indexes.

The B language rules concerning the possible values of the *BAR_VALUE* parameter ensure that a computing variable being able to contain the elements of MININT..MAXINT, then it can contain those of *BAR_VALUE*. For example, the B rules do not permit to give to *BAR_VALUE* the value MAXINT+1,MAXINT+2.

VAL_ARR_RGE

syntax $vv \leftarrow \text{VAL_ARR_RGE}(rr,ii)$

preconditions *ii* must be a *BAR_INDEX*, *rr* must be a *BAR_RANGE*

outputs *vv* is an element of *BAR_VALUE*, which value is the array value at position *ii*, line *rr*.

STR_ARR_RGE

syntax $\text{STR_ARR_RGE}(rr,ii,vv)$

preconditions *rr* must be an element of *BAR_RANGE*, *ii* an element of *BAR_INDEX* and *vv* an element of *BAR_VALUE*

Value *vv* is stored in the array line *rr*, index *ii*.

COP_ARR_RGE

syntax $\text{COP_ARR_RGE}(\text{dest},\text{src})$

preconditions *dest* and *src* must be elements of *BAR_RANGE*

The *src* line is copied to the *dest* line.

CMP_ARR_RGE

syntax $bb \leftarrow \text{CMP_ARR_RGE}(\text{range1},\text{range2})$

preconditions *range1* and *range2* must be elements of *BAR_RANGE*

outputs *bb* is an element of *BOOL*, that takes the TRUE value if the two lines are equal.

C++ LANGUAGE

In C++, the array is realised by an array of pointers, pointing on integers arrays. The access to these arrays are done using methods that refuse the index used between 0 and the arrays size, guaranting an optimal memory occupation.

The memory is dynamically reserved when launching the program. If the size indicated by the formal parameters is too big, the program stops with the following message:

Virtual memory exceeded in 'new'

C LANGUAGE

The realisation in C is based on the same principles as in C++. The stop message on the initial reservation failure is:

```
Fatal error: Malloc of X bytes failed
Execution of current application is aborted
```

ADA LANGUAGE

The use of generic packages guarantees an optimal memory occupation. No restriction is made on the instancing parameters. On an initial reservation failure, an exception stops the program.

4.3 BASIC_IO: vt100 style inputs/outputs

OPERATIONS

INTERVAL_READ operator input of an integer in mm..nn.
 INT_WRITE print an integer.
 BOOL_READ operator input of a Boolean TRUE or FALSE state
 BOOL_WRITE print TRUE or FALSE.
 CHAR_READ operator input of a character.
 CHAR_WRITE print a character.
 STRING_WRITE print a message.

SIMPLE EXAMPLE

The following implementation displays “hello” on the terminal:

<pre> MACHINE bonj OPERATIONS main = skip END </pre>	<pre> IMPLEMENTATION bonj_1 REFINES bonj IMPORTS BASIC_IO OPERATIONS main = BEGIN STRING_WRITE("hello\n") END END </pre>
---	---

DESCRIPTION

BASIC_IO is used for simple input/output actions on a terminal. This basic machine is used to build models. Such I/O cannot be considered as safe.

In UNIX, the system devices used are standard input and standard output (*stdin* and *stdout*), they can therefore be redirected.

INTERVAL_READ

syntax $bb \leftarrow \text{INTERVAL_READ}(mm,nn)$
preconditions mm and nn must be NATs so that $mm \leq nn$
outputs bb integer in $mm..nn$

The operator inputs an integer of the interval $mm..nn$. The input format forces to type a succession of number(s) followed by RETURN. The first input character must be a number. On the opposite case, the input fails “3” is not valid). When a character that is not the first input is not a number anymore, this character, as all the following ones, are ignored: “3e2” is a valid input of the integer 3. As long as the input is false, the message “THIS IS NOT A NUMBER IN $mm..nn$ ” is displayed and a new entry is required.

INT_WRITE

syntax INT_WRITE(vv)

preconditions vv must belong to NAT

Output number vv, with no return.

BOOL_READ

syntax bb ← BOOL_READ

outputs bb must be Boolean.

The operator enters Boolean TRUE or FALSE conditions, with no character before it (for example: “TRUE” is rejected because of the space before it). As long as the operator has not made a valid entry, the message “THIS IS NOT A BOOL VALUE: type TRUE or FALSE” is displayed and a new entry is required.

BOOL_WRITE

syntax BOOL_WRITE(bb)

preconditions bb must be Boolean

Output TRUE or FALSE, with no return.

CHAR_READ

syntax cc ← CHAR_READ

outputs cc must be part of 0..255

Operator entry of a character that is interpreted as a number in 0..255. Type in the character followed by return. If several characters has been typed, only the first one is taken into account (example: “cdef” is understood as “=32). In C, pressing Return only returns 10, ctrl-D (EOF) returns 0. In ADA, only the ‘visible’ characters entries (i.e, no control characters) are accepted.

CHAR_WRITE

syntax CHAR_WRITE(vv)

preconditions vv must belong to the range 0..255

Displays the cc character on-screen (example: CHAR_WRITE(10) to produce a return). Remember, a single quote means “prime” the language’s notation conventions, and B. CHAR_WRITE('A') for example, means nothing. On the contrary, the quoted strings are valid elements in a formula, they serve for STRING_WRITE below.

STRING_WRITE

syntax STRING_WRITE(ss)

preconditions ss must be an element in the STRING set

Will display a character string on-screen. For ss use quoted strings. A “C type” formatting is used, even for a translation into ADA, i.e.:

\t produces a tab

\E produces Escape

\B produces a sound

\" produces a quote

KNOWN PROBLEMS

STRING does not have a coherent definition. The prover proves that any character string belongs to STRING due to an *ad hoc* rule, that does not derive from the definition `STRING = seq(CHAR)`. In addition, using a STRING type local variable in an implementation is not possible. To be completely rigorous, nothing ensures that the operator performs all the requested entries. Therefore the operations for entering the true data entry module (BASIC_IO.c for example) do not really implant the specifications of the corresponding B operations.

PROGRAMMING

A more complete example:

<div><div><div><div><div>MACHINE</div><div>bio</div></div><div><div>OPERATIONS</div><div>main = skip</div></div><div><div>END</div></div></div></div></div>	<div><div><div>IMPLEMENTATION</div><div>bio_1</div></div><div><div>REFINES</div><div>bio</div></div><div><div>IMPORTS</div><div>BASIC_ARITHMETIC,BASIC_IO</div></div><div><div>OPERATIONS</div><div>main = VAR zz,bb,cc IN</div><div>zz ← INTERVAL_READ(0,100);</div><div>STRING_WRITE("this is the value : ");</div><div>INT_WRITE(zz);</div><div>CHAR_WRITE(10);</div><div>bb ← BOOL_READ;</div><div>STRING_WRITE("this is the value : ");</div><div>BOOL_WRITE(bb);</div><div>CHAR_WRITE(10);</div><div>cc ← CHAR_READ;</div><div>STRING_WRITE("this is the value : ");</div><div>INT_WRITE(cc);</div><div>STRING_WRITE(" = ");</div><div>CHAR_WRITE(cc);</div><div>CHAR_WRITE(10)</div><div>END</div><div>END</div></div></div>
---	---

Execution example:

```
ATELIER-B% bio
sdfsdf
THIS IS NOT A NUMBER IN 0..100
20
this is the value: 20
CRUE
THIS IS NOT A BOOL VALUE: type TRUE or FALSE
TRUE
this is the value: TRUE
cvf
```

```
this is the value: 99 = c
ATELIER-B%
```

NOTE: To be completely rigorous, nothing ensures that the operator performs all the entries requested. The entry loops of the concrete module (BASIC_IO.c for example) do not really implant the specifications of the corresponding operations.

Possible evolutions:

It should be possible to define in the machine BASIC_IO., abstract variables modeling the inputs/outputs; it should then be possible to specify the required interactions *of the external system*. The abstract machine that needs to handle inputs/outputs will use BASIC_IO notions (by SEES or INCLUDES) to represent the required interactions.

5 Description of Library Machines

The library machines are all intended for creating mathematical objects, except machine `L_ARITHMETIC1` that provides certain arithmetical functions. The modeled mathematical objects are:

total functions : these are machines contain “ARR” (array) in their name;

partial functions : machines with the “PFNC” (partial function) in their name;

sets : these are machines with the “SET” (set) in their name;

sequences : these are machines with the “SEQ” (sequence) in their name.

For each mathematical object, it is possible to realize either a variable representing the object, or a variable representing several objects of this type. For each type of object, it is therefore possible to realize:

- The object itself;
- An array of objects with the same type, same size, these are machines with a name containing the “RGE” (range) radical;
- A partial function of objects with the same size and same type, these are machines with a name containing the “COL” (collection) radical;
- A partial function of objects with the same type, but with various sizes (“OBJ” radical).

The “RGE” and “COL” type machines produce objects that consume the memory necessary for the maximum number of required objects. For example, if we create a range or a collection of three sequences of at least ten elements, we will always require 30 memory spaces; but the use of a collection avoids the user program to manage the sequences available/occupied. Object machines reserve a memory space that may be freely distributed depending on the created objects and their size. Mathematical objects listed above are not all available on the different types of machines, refer to library machines table of contents for the list that corresponds to the current version.

WARNING : Most of the library machines are based on the basic machines `BASIC_ARRAY_VAR` and `BASIC_ARRAY_RGE`. The manual implementations of the basic machines `BASIC_ARRAY_VAR` and `BASIC_ARRAY_RGE` destined to the translators supplied with Atelier B are provided as a demonstration. They are not safe, and not appropriate in all the B use context. In the case of a more complete use, the user would have to realize these basic machines.

5.1 L_ARITHMETIC1: Extended Integer Operations

The “integer” term refers to the elements of NAT.NAT that is the set of the natural integers between 0 and MAXINT.

OPERATIONS

MIN	minimum of two integers.
MAX	maximum of two integers.
INC	increment an integer strictly inferior to MAXINT.
DEC	decrement a literal integer.
EXP	exponentiation.
SQRT	default integer square root.
LOG_BY_DEFAULT	default logarithm.
LOG_BY_EXCESS	logarithm by excess.

EXAMPLE

The example below shows a machine that uses a certain number of functionalities of the machine L_ARITHMETIC1.

MACHINE	IMPLEMENTATION
<pre> m1 OPERATIONS xx ← op1 = ANY tt WHERE tt ∈ NAT ∧ tt×tt = 16 THEN xx:=tt END; xx ← op2 = ANY tt WHERE tt ∈ NAT ∧ 3^{tt} = 27 THEN xx:=tt END END </pre>	<pre> m1_1 REFINES m1 IMPORTS L_ARITHMETIC1, OPERATIONS xx ← op1 = BEGIN xx ← SQRT(16) END; xx ← op2 = VAR rr IN xx,rr ← LOG_BY_DEFAULT (3, 27) END END </pre>

DESCRIPTION

L_ARITHMETIC1 offers arithmetical operations such as roots and logarithms, operations on the elements NAT and dedicated to calculatory applications. Calculus being integers values, the search operation for the logarithm and the square root return the *best*¹ approaching value in NAT. The used algorithms are optimized.

MACHINE PARAMETERS

None.

¹ The NAT element immediatly inferior or superior wether the calcul is performed by inferior value or superior value

MIN

syntax $uu \leftarrow \text{MIN}(vv, ww)$
preconditions vv and ww must be in NAT.
outputs $uu = \min(\{vv, ww\})$

MAX

syntax $uu \leftarrow \text{MAX}(vv, ww)$
preconditions vv and ww must be in NAT.
outputs uu receives $\max(\{vv, ww\})$

INC

syntax $uu \leftarrow \text{INC}(vv)$
preconditions vv must be in $0..MAXINT-1$.
outputs $uu = vv + 1$

DEC

syntax $uu \leftarrow \text{DEC}(vv)$
preconditions vv must be in $1..MAXINT$.
outputs $uu = vv - 1$

EXP

syntax $rr \leftarrow \text{EXP}(xx, nn)$
preconditions xx and nn must be in NAT. xx and nn must not both be nil. xx^{nn} must be less than or equal to $MAXINT$.
outputs rr receives xx^{nn}

EXP returns xx to the power of nn . Calculating 0^0 is illegal (0^0 is not defined). The implementation uses a fast algorithm based on breaking down into base 2 of nn ($\log_2(nn)$ iterations).

SQRT

syntax $nn \leftarrow \text{SQRT}(pp)$
preconditions pp must be in NAT.
outputs nn so that $nn \times nn \leq pp < (nn+1) \times (nn+1)$

SQRT returns the largest nn so that $nn \times nn \leq pp$. The implementation uses an algorithm that performs $\text{SQRT}(nn)$ iterations, where each iteration costs two additions and a subtraction.

LOG_BY_DEFAULT

syntax $uu, rr \leftarrow \text{LOG_BY_DEFAULT}(vv, ww)$
preconditions ww and vv are two natural integers and vv is between 2 and $MAXINT$.
outputs uu is the smallest natural so that $vv^{(uu+1)}$ is strictly greater than ww . By definition, uu is a natural integer. rr takes the value vv^{uu} .

`LOG_BY_DEFAULT` in base `vv` of `ww`: returns the smallest `uu` value so that $ww < vv^{(uu+1)}$. This gives $vv^{uu} \leq ww$, except if $ww < vv$ (example: $ww = 0$). Does not work for $vv = 0$ or 1 as 0^{ii} and 1^{ii} are constants. `rr` receives the value of vv^{uu} , which easily allows judging the error made.

LOG_BY_EXCESS

syntax $uu, bb \leftarrow \text{LOG_BY_EXCESS}(vv, ww)$

preconditions `ww` belongs to `NAT` and `vv` is an element of the interval `2..MAXINT`.

outputs `uu` receives the smallest natural so that vv^{uu} is greater than or equal to `ww`. `uu` must be in `NAT`. `bb` is an element of `BOOL`, it indicates whether the logarithm is an exact one.

`LOG_BY_EXCESS` in base `vv` in `ww`: returns the smallest `uu` so that $ww \leq vv^{uu}$. WARNING: vv^{uu} may exceed `MAXINT`! Does not work for $vv = 0$ or 1 as 0^{ii} and 1^{ii} are constants. `bb` equals `TRUE` if $ww = vv^{uu}$.

IMPORTS REQUIRED

None.

5.2 L_ARRAY1: One Dimensional Array, with Initialization Loop

OPERATIONS

VAL_ARRAY value of an element (promoted operation)
 STR_ARRAY write an element (promoted operation)
 SET_ARRAY write the same value in a portion of the array

EXAMPLE

Use SET_ARRAY to initialize an array:

MACHINE m1 VARIABLES vv INVARIANT $vv \in 0..10 \rightarrow 0..255$ INITIALISATION $vv := (0..10) \times \{5\}$ END	IMPLEMENTATION m1_1 REFINES m1 IMPORTS i1.L_ARRAY1(0..255,10) INVARIANT <i>(arr_vrb is the variable in L_ARRAY1)</i> $i1.arr_vrb = vv$ INITIALISATION $i1.SET_ARRAY(0,10,5)$ END
--	---

DESCRIPTION

As it is possible, L_ARRAY1 is used instead of BASIC_ARRAY_VAR. L_ARRAY1 realises, using an array, an abstract variable representing a function. It is then possible to have an initialization operation of the entire function or of a part of it (initialization loop). The starting part of the function performed is an interval: if not, it would not be possible to indicate a portion of this set without mentioning all elements involved.

MACHINE PARAMETERS

L_ARRAY1(LAU_VALUE, LAU_maxidx): LAU_VALUE is the set of possible values for the array elements, 0..LAU_maxidx is the set of array indexes.

VAL_ARRAY

syntax $vv \leftarrow VAL_ARRAY(ii)$
preconditions ii must be in $0..LAU_maxidx$
outputs vv is an element of LAU_VALUE, the array value at position ii .

STR_ARRAY

syntax STR_ARRAY(ii,vv)

preconditions ii and vv must belong to the 0..LAU_maxidx and LAU_VALUE respectively.

vv value is stored in the array at ii index. **SET_ARRAY**

syntax SET_ARRAY (ii,jj,vv)

preconditions ii..jj is a sub-set of 0..LAU_maxidx and vv an element of LAU_VALUE. For implementation reasons ², jj and MAXINT must be different.

The value vv is stored in the array for all the indexes between ii to jj. If ii>jj, the array does not change.

Note that it would not have been advisable to set $ii \leq jj$ as a precondition of this operation, as this would have limited its use. Let us consider the case of a call to SET_ARRAY in a loop. The last iteration of the loop contains a call with the form SET_ARRAY (ii, jj, vv) with $ii = jj + 1$. The presence of a precondition in the definition of the operation SET_ARRAY would force us to “guard” all the calls to SET_ARRAY by an IF. More generally, the precondition must be selected as minimal to protect us from a B code of “*defensive*” aspect.

IMPORTS REQUIRED

None.

WARNING: The implementation of this machine creates the default instance for the BASIC_ARRAY_VAR machine (IMPORTS BASIC_ARRAY_VAR(...)). The addition of an instance of the machine BASIC_ARRAY_VAR requires choosing a new instance name, as, for example: i1.BASIC_ARRAY_VAR).

²Indeed, the loops used make a pre-incrementation, that does not produce literal excedent)

5.3 L_ARRAY3: Array with Non Ordered Values, Maximum Operations

OPERATIONS

- VAL_ARRAYvalue of an element (promoted operation).
- STR_ARRAYwrite an element (promoted operation).
- SET_ARRAYwrite the same value in an array portion (promoted operation).
- SWAP_ARRAYexchange two elements (promoted operation).
- RIGHT_SHIFT_ARRAYshift a portion to the main index (promoted operation).
- LEFT_SHIFT_ARRAYshift a portion to the small index (promoted operation).
- SEARCH_MAX_EQL_ARRAYsearch for a value in an array (promoted operation).
- SEARCH_MIN_EQL_ARRAYsearch for a value in an array portion (promoted operation).
- REVERSE_ARRAYreverse the order of elements in an array portion.

EXAMPLE

The example below is a machine that represents the color assigned to 101 points, this color may be red, green or blue for each point. An operation is used to find a red dot.

MACHINE m1 SETS COLOR = {red, green, blue} VARIABLES color INVARIANT color ∈ 0..100 → COLOR INITIALISATION color :=(0..100) × {red} OPERATIONS ii,bb ← trouve_red = PRE rouge ∈ ran(color) THEN ii :∈ color ⁻¹ [{red}] bb :∈ BOOL END END	IMPLEMENTATION m1_l REFINES m1 IMPORTS i1.L_ARRAY3(COLOR,100) INVARIANT i1.arr_vrb = color INITIALISATION i1.SET_ARRAY(0,100,red) OPERATIONS ii,bb ← trouve_red = VAR bb IN ii,bb ← i1.SEARCH_MAX_EQL_ARRAY(0,100,red) END END
---	---

DESCRIPTION

L_ARRAY3 is the most complete of the one dimensional array machines that do not require that the output set be part of an interval. L_ARRAY5 has been constrained. It is therefore possible to create arrays with values that are elements of a listed set while having access to complete operations such as element order reversal. The operation that

is not available is the one that would require an order relationship on the array elements: sort.

MACHINE PARAMETERS

L_ARRAY3(LAT_VALUE,LAT_maxidx): LAT_VALUE is the set of possible values for array elements, 0..LAT_maxidx is the set of array indexes.

VAL_ARRAY

syntax $vv \leftarrow \text{VAL_ARRAY}(ii)$
preconditions ii must be in 0..LAT_maxidx
outputs vv is a LAT_VALUE, it is the value of the array at position ii .

STR_ARRAY

syntax **STR_ARRAY**(ii,vv)
preconditions ii must be in 0..LAT_maxidx and vv must belong to LAT_VALUE

The vv value is stored in the array at index ii .

SET_ARRAY

syntax **SET_ARRAY**(ii,jj,vv)
preconditions $ii..jj$ must be a subset of 0..LAT_maxidx and vv belong to LAT_VALUE.
 For implementation reasons it is also necessary that jj be different from MAXINT.

The vv value is stored in the array for all indexes between ii and jj . If $ii > jj$, the array will not change.

SWAP_ARRAY

syntax **SWAP_ARRAY**(ii,jj)
preconditions ii,jj must be in 0..LAT_maxidx.

The ii and jj elements in the array are exchanged.

RIGHT_SHIFT_ARRAY

syntax **RIGHT_SHIFT_ARRAY**(ii,jj,nn)
preconditions ii,jj,nn must be in 0..LAT_maxidx, with $ii \leq jj$ and $jj+nn \leq \text{LAT_maxidx}$
 to make possible the possible the shift to the right by nn spaces.

Part $ii+nn..jj+nn$ receives a copy of part $ii..jj$ of the array (shift nn spaces to the right).

LEFT_SHIFT_ARRAY

syntax **LEFT_SHIFT_ARRAY**(ii,jj,nn)
preconditions ii,jj must be in 0..LAT_maxidx, with $ii \leq jj$. nn must be NAT with $nn \leq ii$
 to make possible the shift to the left by nn places. For implementation reasons, jj must be not equal MAXINT.

The $ii-nn..jj-nn$ part receives a copy of part $ii..jj$ from the array (shift nn spaces to the left).

SEARCH_MAX_EQL_ARRAY

syntax $rr, bb \leftarrow \text{SEARCH_MAX_EQL_ARRAY}(ii, jj, vv)$

preconditions ii and jj must be in $0..LAT_maxidx$, $ii \leq jj$ and vv belong to LAT_VALUE .

outputs TRUE if vv was found, FALSE if not. rr is a NAT, if $bb = \text{TRUE}$ then rr is the largest index in the array worth vv .

Search for an array element equal to vv , by scanning the $ii..jj$ part starting from jj .

SEARCH_MIN_EQL_ARRAY

syntax $rr, bb \leftarrow \text{SEARCH_MIN_EQL_ARRAY}(ii, jj, vv)$

preconditions ii and jj must be in $0..LAT_maxidx$, $ii \leq jj$ and vv belong to LAT_VALUE .

outputs TRUE if vv was found, FALSE if not. rr is a NAT, if $bb = \text{TRUE}$, then rr is the smallest index in the array worth vv .

Search for an array element that equals vv , by scanning the $ii..jj$ part starting from ii .

REVERSE_ARRAY

syntax $\text{REVERSE_ARRAY}(ii, jj)$

preconditions ii and jj must be in $0..LAT_maxidx$.

Reverse the order of elements in the $ii..jj$ portion of the array.

IMPORTS REQUIRED

(instances to import as the implementation tree for this library machine sees them with SEES)

BASIC_ARITHMETIC; BASIC_BOOL.

WARNING: The implementation of this machine creates the default instance for the BASIC_ARRAY_VAR machine (clause IMPORTS BASIC_ARRAY_VAR(...)). Therefore if another instance is necessary, it must be given a different instance name (for example: i1.BASIC_ARRAY_VAR).

5.4 L_ARRAY5: Array with Ordered Values, Sort Operation

OPERATIONS

VAL_ARRAY	value of an element (promoted operation).
STR_ARRAY	write an element (promoted operation).
SET_ARRAY	write the same value to a portion of an array (promoted operation).
SWAP_ARRAY	exchange two elements (promoted operation).
RIGHT_SHIFT_ARRAY	shift a portion to the large index (promoted operation).
LEFT_SHIFT_ARRAY	shift a portion to the small index (promoted operation).
SEARCH_MAX_EQL_ARRAY	search for a value in a portion of the array (promoted operation).
SEARCH_MIN_EQL_ARRAY	search for a value in a portion of the array (promoted operation).
REVERSE_ARRAY	reverse the order of the elements in a portion of the array (promoted operation).
SEARCH_MIN_GEQ_ARRAY	search for the first element that exceeds a value (promoted operation).
ASCENDING_SORT_ARRAY	sort a portion of the array.

EXAMPLE

MACHINE m1 VARIABLES vv INVARIANT $vv \in 0..4 \rightarrow 0..255 \wedge$ $\forall xx.(xx \in 0..3 \Rightarrow$ $vv(xx) \geq vv(xx+1))$ INITIALISATION $vv : (vv \in 0..4 \rightarrow 0..255 \wedge$ $\forall xx.(xx \in 0..3 \Rightarrow$ $vv(xx) \geq vv(xx+1)))$ END	IMPLEMENTATION m1_1 REFINES m1 IMPORTS L_ARRAY5(0,255,4) INVARIANT arr_vrb = vv INITIALISATION SET_ARRAY(0,4,50); STR_ARRAY(2,10); STR_ARRAY(4,30); ASCENDING_SORT_ARRAY(0,4); REVERSE_ARRAY(0,4) END
--	---

DESCRIPTION

L_ARRAY5 is the most complete of the one dimensional array machines. It especially comprises a sort operation implanted using a shift sort (fast algorithm).

MACHINE PARAMETERS

$L_ARRAY5(LAC_minval, LAC_maxval, LAC_maxidx)$: $LAC_minval..LAC_maxval$ is the set of possible values for the elements in the array, $0..LAC_maxidx$ is the set of index values for the array. LAC_minval , LAC_maxval , LAC_maxidx must be NATs: this machine does not allow negative values. It is also necessary for $LAC_minval \leq LAC_maxval$ and $1 \leq LAC_maxidx$.

VAL_ARRAY

syntax $vv \leftarrow VAL_ARRAY(ii)$
preconditions ii must be in $0..LAC_maxidx$
outputs vv is in $LAC_minval..LAC_maxval$, is the array value at position ii .

STR_ARRAY

syntax $STR_ARRAY(ii, vv)$
preconditions ii must be in $0..LAC_maxidx$ and vv in $LAC_minval..LAC_maxval$ and LAC_VALUE .

The vv value is stored in the array at index ii .

SET_ARRAY

syntax $SET_ARRAY(ii, jj, vv)$
preconditions $ii..jj$ must be included in $0..LAC_maxidx$ and vv must be in LAC_VALUE . For implementation, it is also necessary that jj be different from the $MAXINT$ constant.

The vv value is stored in the array for all indexes from ii to jj . If $ii > jj$, the array does not change.

SWAP_ARRAY

syntax $SWAP_ARRAY(ii, jj)$
preconditions ii, jj must be in $0..LAC_maxidx$.

The ii and jj elements in the array are exchanged.

RIGHT_SHIFT_ARRAY

syntax $RIGHT_SHIFT_ARRAY(ii, jj, nn)$
preconditions ii, jj, nn must be in $0..LAC_maxidx$, with $ii \leq jj$ and $jj + nn \leq LAC_maxidx$ to make possible the right shift by nn spaces.

The $ii + nn..jj + nn$ part receives a copy of the $ii..jj$ part of the array (shift right by nn spaces).

LEFT_SHIFT_ARRAY

syntax $LEFT_SHIFT_ARRAY(ii, jj, nn)$
preconditions ii, jj must be in $0..LAC_maxidx$, with $ii \leq jj$. nn must be a NAT with $nn \leq ii$ to allow the left shift by nn spaces. For implementation reasons, jj cannot equal $MAXINT$.

The $ii - nn..jj - nn$ part receives a copy of the $ii..jj$ part of the array (shift left by nn spaces).

SEARCH_MAX_EQL_ARRAY

syntax $rr, bb \leftarrow \text{SEARCH_MAX_EQL_ARRAY}(ii, jj, vv)$

preconditions ii and jj must be in $0..LAC_maxidx$, $ii \leq jj$ and vv be in LAC_VALUE .

outputs $TRUE$ if vv was found, $FALSE$ if not. rr is a NAT, if $bb = TRUE$, then rr is the highest index in the array worth vv .

Search for an array element equal to vv , by scanning the $ii..jj$ part starting from jj .

SEARCH_MIN_EQL_ARRAY

syntax $rr, bb \leftarrow \text{SEARCH_MIN_EQL_ARRAY}(ii, jj, vv)$

preconditions ii and jj must be in $0..LAC_maxidx$, $ii \leq jj$ and vv be in LAC_VALUE .

outputs $TRUE$ if vv was found, $FALSE$ if not. rr is a NAT, if $bb = TRUE$ then rr is the smallest index in the array worth vv .

Search for an array element equal to vv , by scanning the $ii..jj$ part starting from ii .

REVERSE_ARRAY

syntax $\text{REVERSE_ARRAY}(ii, jj)$

preconditions ii and jj must be in $0..LAC_maxidx$.

Reverse the order of the elements in the $ii..jj$ portion of the array.

SEARCH_MIN_GEQ_ARRAY

syntax $ii, bb \leftarrow \text{SEARCH_MIN_GEQ_ARRAY}(jj, kk, vv)$

preconditions jj and kk must be in $0..LAC_maxidx$, $jj \leq kk$ and vv be in $LAC_minval..LAC_maxval$. For implementation location reasons, kk must not equal the $MAXINT$ constant.

outputs $TRUE$ if an element that is greater or equal to vv was found, $FALSE$ if not. ii is a NAT, if $bb = TRUE$, then ii is the smallest index in the image array that is greater than or equal to vv .

Search for an element that is greater than or equal to vv in $jj..kk$ starting from jj .

ASCENDING_SORT_ARRAY

syntax $\text{ASCENDING_SORT_ARRAY}(ii, jj)$

preconditions ii and jj must be in $0..LAC_maxidx$. For implementation reasons, ii and jj must not equal $MAXINT$.

Shift sort, in ascending order (the smallest first) on the $ii..jj$ portion.

IMPORTS REQUIRED

(instances to import as the implementation tree for this library machine sees them with SEES)

BASIC_ARITHMETIC; BASIC_BOOL.

WARNING: The implementation of this machine creates the default instance for the BASIC_ARRAY_VAR machine (clause IMPORTS BASIC_ARRAY_VAR(...)). Therefore if another instance is required it must be given a different instance name (for example: i1.BASIC_ARRAY_VAR).

5.5 L_PFNC: Partial Function

OPERATIONS

VAL_PFNC	value of the function for an element in its domain
STR_PFNC	overloads the partial function with a pair
XST_PFNC	test that an index is in the partial function domain
RMV_PFNC	removes a pair from the partial function
SET_PFNC	overloads a part of the function with a constant
SWAP_PFNC	exchanges the images for two domain indexes
RIGHT_SHIFT_PFNC	right shift part of the domain
LEFT_SHIFT_PFNC	left shift part of the domain
SEARCH_MAX_EQL_PFNC	search for a value in the partial function
SEARCH_MIN_EQL_PFNC	search for a value in the partial function
REVERSE_PFNC	reverse the order of elements in a portion of the domain
ASCENDING_SORT_PFNC	sort in a portion of the domain

EXAMPLE

MACHINE m1 VARIABLES pf INVARIANT $pf \in 0..10 \leftrightarrow 0..255$ INITIALISATION $pf := \{4 \mapsto 6\}$ END	IMPLEMENTATION m1_1 REFINES m1 IMPORTS L_PFNC(0,255,10) INVARIANT pfnc_vrb = pf INITIALISATION STR_PFNC(4,6) END
---	--

DESCRIPTION

L_PFNC implements a partial function with almost all of the operations available in L_ARRAY5 (In fact only SEARCH_MIN_GEQ is not used). The practical usefulness of partial functions is that they dispense with the need to add a "non existent" or "unused" element in the input sets in order to implant them as total functions. The implementation of L_PFNC performs these elements by using the seldom used MAXINT value.

MACHINE PARAMETERS

L_PFNC(LPF_minval,LPF_maxval,LPF_maxidx): LPF_minval..LPF_maxval is the input set of the function, 0..LPF_maxidx is the source set. LPF_minval, LPF_maxval,

LPF_maxidx must be NATs: this machine does not allow negative values. Moreover, $\text{LPF_minval} \leq \text{LPF_maxval}$ and $1 \leq \text{LPF_maxidx}$; as well as $\text{LPF_maxval} < \text{MAXINT}$: This is because MAXINT is used to indicate that the corresponding index is not part of the partial function. Again to simplify implementation, it is also illegal to have $\text{LPF_maxidx} = \text{MAXINT}$.

VAL_PFNC

syntax $\text{vv} \leftarrow \text{VAL_PFNC}(\text{ii})$
preconditions ii must be in the partial function domain
outputs vv is in $\text{LPF_minval}..\text{LPF_maxval}$, it is the value of the array at position ii.

STR_PFNC

syntax $\text{STR_PFNC}(\text{ii}, \text{vv})$
preconditions ii must be in $0..\text{LPF_maxidx}$ and vv be in $\text{LPF_minval}..\text{LPF_maxval}$.

The partial function is overloaded by $\{\text{ii} \mapsto \text{vv}\}$.

XST_PFNC

syntax $\text{bb} \leftarrow \text{XST_PFNC}(\text{ii})$
outputs bb is TRUE if ii is in the domain of the function, FALSE if not.

RMV_PFNC

syntax $\text{RMV_PFNC}(\text{ii})$
preconditions ii must be in the domain of the partial function.

The $\{\text{ii} \mapsto \text{pfnc_vrb}(\text{ii})\}$ pair is removed from the partial function pfnc_vrb.

SET_PFNC

syntax $\text{SET_PFNC}(\text{ii}, \text{jj}, \text{vv})$
preconditions ii..jj must be included in $0..\text{LPF_maxidx}$ and vv be in $\text{LPF_minval}..\text{LPF_maxval}$. ii and jj must be NATs.

The partial function is overloaded by $(\text{ii}..\text{jj}) \times \text{vv}$. If $\text{ii} > \text{jj}$,

ii..jj is blank and the partial function is not modified, but it is still necessary for ii and jj to be NATs.

SWAP_PFNC

syntax $\text{SWAP_PFNC}(\text{ii}, \text{jj})$
preconditions ii,jj must be in the domain of the partial function.

The ii and jj elements in the array are exchanged.

RIGHT_SHIFT_PFNC

syntax **RIGHT_SHIFT_PFNC**(ii,jj,nn)

preconditions ii,jj,nn must be in 0..LPF_maxidx, with $ii \leq jj$ and $jj+nn \leq LPF_maxidx$ to allow the right shift by nn spaces. It is also necessary for ii..jj to be included in the domain of the partial function.

The ii+nn..jj+nn part is overloaded by a copy of the ii..jj part in the partial function (shift by nn spaces to the right).

LEFT_SHIFT_PFNC

syntax **LEFT_SHIFT_PFNC**(ii,jj,nn)

preconditions ii,jj must be in 0..LPF_maxidx, with $ii \leq jj$. nn must be a NAT with $nn \leq ii$ to allow the left shift by nn spaces. In addition it is necessary for ii..jj to be included in the domain of the partial function.

The ii-nn..jj-nn part is overloaded by a copy of the ii..jj part in the partial function (shift left by nn spaces).

SEARCH_MAX_EQL_PFNC

syntax rr,bb \leftarrow **SEARCH_MAX_EQL_PFNC**(ii,jj,vv)

preconditions ii and jj must be in 0..LPF_maxidx, $ii \leq jj$ and vv be in LPF_minval..LPF_maxval.

outputs TRUE if vv was found, FALSE if not, rr is a NAT, if bb = TRUE, then rr is the largest index, the image of which by the partial function is vv.

Search for an array element that equals vv, by scanning the ii..jj part, starting from jj.

SEARCH_MIN_EQL_PFNC

syntax rr,bb \leftarrow **SEARCH_MIN_EQL_PFNC**(ii,jj,vv)

preconditions ii and jj must be in 0..LPF_maxidx, $ii \leq jj$ and vv be in LPF_minval..LPF_maxval.

outputs TRUE if vv was found, FALSE if not, rr is a NAT, if bb = TRUE, then rr is the smallest index, the image of which by the partial function is vv.

Search for an array element that equals vv, by scanning the ii..jj part starting from ii.

REVERSE_PFNC

syntax **REVERSE_PFNC**(ii,jj)

preconditions ii and jj must be in 0..LPF_maxidx, and ii..jj must be included in the domain of the partial function.

Reverse the order of the elements in the ii..jj portion of the partial function.

ASCENDING_SORT_PFNC

syntax **ASCENDING_SORT_PFNC**(ii,jj)

preconditions ii and jj must be in 0..LPF_maxidx, and ii..jj must be included in the domain of the partial function.

Shift sort, in ascending order (the smallest first) in the ii..jj portion.

IMPORTS REQUIRED

(instances to import as the implementation tree for this library machine sees them with SEES) BASIC_ARITHMETIC; BASIC_BOOL.

WARNING: The implementation of this machine creates the default instance for the BASIC_ARRAY_VAR machine (clause IMPORTS BASIC_ARRAY_VAR(...)). Therefore if another instance is necessary, it must be given a different instance name (for example: i1.BASIC_ARRAY_VAR).

5.6 L_SEQUENCE: Creating a Sequence

OPERATIONS

- LEN_SEQ returns the current size of the sequence.
- IS_FULL_SEQ shows whether the sequence is full (size = LS_maxsize).
- IS_INDEX_SEQ shows whether ii is a valid index.
- VAL_SEQ value of an element in the sequence.
- FIRST_SEQ returns the first element in the sequence.
- LAST_SEQ returns the last element in the sequence.
- PUSH_SEQ adds vv to the end of the sequence.
- POP_SEQ removes the last element from the sequence (its value is lost).
- STR_SEQ changes the value of an element in the sequence.
- RMV_SEQ removes an element from the middle of the sequence.
- INS_AFT_SEQ inserts vv right after index ii.
- CLR_SEQ clears the sequence.
- TAIL_SEQ removes the first element from the sequence.
- KEEP_SEQ only keeps the nn first elements in the sequence.
- CUT_SEQ cuts the nn first elements from the sequence.
- PART_SEQ only keeps the ii..jj portion in the sequence.
- REV_SEQ reverses the order of the elements in the sequence.
- FIND_FIRST_SEQ searches for vv in the sequence, starting from the beginning.
- FIND_LAST_SEQ searches for vv in the sequence, starting from the end.

EXAMPLE

The example below shows the use of L_SEQUENCE for a listed set.

MACHINE m1 SETS ST = {classic,baroque} VARIABLES vv INVARIANT vv ∈ seq(ST) ∧ size(vv) ≤ 10 INITIALISATION vv := [baroque,baroque]	IMPLEMENTATION m1_l REFINES m1 IMPORTS L_SEQUENCE(10,ST) INVARIANT (<i>seq_vrb is the variable in L_SEQUENCE</i>) seq_vrb = vv INITIALISATION PUSH_SEQ(baroque) (<i>L_SEQUENCE guarantees</i> PUSH_SEQ(baroque) <i>that the sequence is empty at the start</i>) END
--	--

DESCRIPTION

L_SEQUENCE provides a sequence type variable, the maximum size of which is a machine parameter. Conventional search and shift functions are provided for the practical use of this sequence. This answers the frequent problem in programming applications which is to maintain a list with no blanks.

MACHINE PARAMETERS

L_SEQUENCE(LS_maxsize,LS_VALUE): the variable is a sequence of LS_VALUE elements, with a maximum size that is LS_maxsize.

LEN_SEQ

syntax $nn \leftarrow \text{LEN_SEQ}$
outputs $0..LS_maxsize$

Returns the current size of the sequence.

IS_FULL_SEQ

syntax $bb \leftarrow \text{IS_FULL_SEQ}$
outputs bb is TRUE if the sequence is full, FALSE if not.

Specifies whether the sequence is full (size = LS_maxsize).

IS_INDEX_SEQ

syntax $bb \leftarrow \text{IS_INDEX_SEQ}(ii)$
preconditions ii must be a NAT.
outputs bb is TRUE if ii is an index in the sequence, FALSE if not.

Specifies whether ii is a valid index.

VAL_SEQ

syntax $vv \leftarrow \text{VAL_SEQ}(ii)$
preconditions ii must be an index in the sequence ($ii \in 1..size(seq_vrb)$).
outputs vv is the value of the ii -ith element ($vv \in \text{VALUE}$).

Value of an element in the sequence.

FIRST_SEQ

syntax $vv \leftarrow \text{FIRST_SEQ}$
preconditions the sequence must not be empty.
outputs vv is the value of the first element ($vv \in \text{VALUE}$).

Returns the first element in the sequence.

LAST_SEQ

syntax $vv \leftarrow \text{LAST_SEQ}$
preconditions the sequence must not be empty.
outputs vv is the value of the last element ($vv \in \text{VALUE}$).

Returns the last element in the sequence.

PUSH_SEQ

syntax PUSH_SEQ(vv)

preconditions vv must be in VALUE and the sequence must not be full.

Add vv at the end of the sequence.

POP_SEQ

syntax POP_SEQ

preconditions the sequence must not be empty.

Removes the last element from the sequence (its value is lost).

STR_SEQ

syntax STR_SEQ(ii,vv)

preconditions vv must be in VALUE and ii must be a valid index for the sequence.

Changes the value of an existing element in the sequence.

RMV_SEQ

syntax RMV_SEQ(ii)

preconditions ii must be a valid index in the sequence.

Removes an element from the middle of the sequence.

INS_AFT_SEQ

syntax INS_AFT_SEQ(ii,vv)

preconditions vv must be in VALUE and ii must be a valid index for the sequence. The sequence must not be full.

Inserts vv right after index ii.

CLR_SEQ

syntax CLR_SEQ

Clears the sequence.

TAIL_SEQ

syntax TAIL_SEQ

preconditions the sequence must not be empty.

Removes the first element from the sequence.

KEEP_SEQ

syntax KEEP_SEQ(nn)

preconditions nn must be a NAT.

Only retains the nn first elements in the sequence. For nn = size(seq_vrb), this operation does not take action.

CUT_SEQ

syntax CUT_SEQ(nn)

preconditions nn must be a NAT.

Deletes the nn first elements from the sequence. For $nn = \text{size}(\text{seq_vrb})$, this operation is equivalent to CLR_SEQ.

PART_SEQ

syntax PART_SEQ(ii,jj)

preconditions ii and jj must be non null NATs, with $ii \leq jj$.

Only retains the ii..jj portion in the sequence. ii..jj may not be included in the domain of the sequence.

REV_SEQ

syntax REV_SEQ

Reverses the order of the elements in the sequence. Applies even for sequences that are empty or of size 1.

FIND_FIRST_SEQ

syntax bb,ii \leftarrow FIND_FIRST_SEQ(vv)

preconditions vv must be in VALUE.

outputs bb is TRUE if vv is in the sequence, FALSE if not. ii belongs to the range 1..LS_maxsize, if bb = TRUE, then it indicates the first position equal to vv.

Search for vv in the sequence, starting from the start.

FIND_LAST_SEQ

syntax bb,ii \leftarrow FIND_LAST_SEQ(vv)

preconditions vv must be in VALUE.

outputs bb is TRUE if vv is in the sequence, FALSE if not. If bb = TRUE, ii belongs to the range 1..LS_maxsize and indicates the last position equal to vv.

Search for vv in the sequence, starting from the end.

IMPORTS REQUIRED

(instances to import as the implementation tree for this library machine sees them with SEES) BASIC_ARITHMETIC; BASIC_BOOL.

WARNING: The implementation of this machine creates the default instance for the BASIC_ARRAY_VAR machine (clause IMPORTS BASIC_ARRAY_VAR(...)). Therefore if another instance is required, it must be given a non blank instance name (for example: i1.BASIC_ARRAY_VAR).

5.7 L_SET: Creating a Set

OPERATIONS

CARD_SET	returns the cardinal for the set.
IS_FULL_SET	identifies whether the set is full (card = LSET_maxsize).
FIND_SET	finds an element in the set.
RMV_SET	removes an element from the set.
INS_SET	inserts an element in the set.
CLR_SET	removes all of the elements from the set.
IS_INDEX_SET	identifies whether a number is a valid index.
VAL_SET	value of an element in the set.

EXAMPLE

The example below shows the use of L_SET on a listed set.

MACHINE	IMPLEMENTATION
m1	m1_1
SETS	REFINES
ST = {cat, dog, bird}	m1
VARIABLES	IMPORTS
vv	L_SET(3,ST)
INVARIANT	INVARIANT
vv \subseteq ST	<i>(set_vrb is the variable in L_SET)</i>
INITIALISATION	ran (set_vrb) = vv
vv := {cat,bird}	INITIALISATION
END	<i>(L_SET ensures that the set is empty at the start)</i>
	INS_SET(cat);
	INS_SET(bird)
	END

DESCRIPTION

L_SET creates a set that is modeled by an injective sequence type variable, set_vrb the maximum size of which is a machine parameter. It offers functions to search for, add and delete elements.

The use of an injective sequence type variable enables easy access to each element of the set via an index. The user can therefore create loops by using the CARD_SET and VAL_SET functions. This would not have been possible if the variable directly represented the set.

WARNING: The user must add the gluing invariant $\text{ran}(\text{set_vrb}) = \text{var_locale}$ to his machine in order to link his set variable with the L_SET machine state.

MACHINE PARAMETERS

L_SET (LSET_maxsize, LSET_VALUE): the variable is an injective sequence of elements from LSET_VALUE, with a maximum size LSET_maxsize.

CARD_SET

syntax $nn \leftarrow \text{CARD_SET}$
output nn is the size of the set (the cardinal of $\text{ran}(\text{set_vrb})$). Therefore, nn belongs to $0.. \text{LSET_maxsize}$

Returns the size of the set.

IS_FULL_SET

syntax $bb \leftarrow \text{IS_FULL_SET}$
output bb is TRUE if the set is full, FALSE if not.

States whether the set is full (size = LSET_maxsize).

IS_INDEX_SET

syntax $bb \leftarrow \text{IS_INDEX_SET}(ii)$
preconditions ii must be a NAT.
outputs bb is TRUE if ii is an index of the set, FALSE if not.

States whether ii is a valid index.

VAL_SET

syntax $vv \leftarrow \text{VAL_SET}(ii)$
preconditions ii must be an index of the set ($ii \in 1..\text{size}(\text{seq_vrb})$).
outputs vv is the value of the ii -the element ($vv \in \text{LSET_VALUE}$).

Value of an element of the set.

FIND_SET

syntax $bb, ii \leftarrow \text{FIND_SET}(vv)$
preconditions vv must be in LSET_VALUE.
outputs bb is TRUE if vv is in the set, FALSE if not. ii is a NAT, if $bb = \text{TRUE}$, then it indicates the position of element vv .

Search for vv in the set.

RMV_SET

syntax $\text{RMV_SET}(vv)$
preconditions vv must be in the set.

Removes an element from the set.

INS_SET

syntax INS_SET(vv)

preconditions vv must be in LSET_VALUE.

Adds an element to the end of the set, if it is not already in it, if not it does nothing.

CLR_SET

syntax CLR_SET

Clears the set.

5.8 L_ARRAY1_RANGE: A Range of Arrays of the Same Size, with Numerical Indexes

OPERATIONS

VAL_ARR_RGE	value of an element (promoted operation).
STR_ARR_RGE	write an element (promoted operation).
COP_ARR_RGE	copy an array to another (promoted operation).
CMP_ARR_RGE	compare two arrays (promoted operation).
DUP_ARR_RGE	duplicate the same array to a series of arrays.
SET_ARR_RGE	copy the same value to an index interval in one of the arrays.
PCOP_ARR_RGE	copy part of one array to a different array, to a given position.
PCMP_ARR_RGE	search for the first element that is different between two parts of two arrays. A Boolean element indicates whether this element was found and, in this case, the index of this element is returned.

EXAMPLE

Using SET_ARR_RGE and DUP_ARR_RGE to initialize a set of arrays:

MACHINE m1 VARIABLES vv INVARIANT $vv \in 0..20 \rightarrow (0..10 \rightarrow 0..255)$ INITIALISATION $vv := (0..20) \times \{(0..10) \times \{5\}\}$ END	IMPLEMENTATION m1_1 REFINES m1 IMPORTS i1.L_ARRAY1_RANGE(0,20,10,0..255) INVARIANT i1.arr_rge = vv INITIALISATION i1.SET_ARR_RGE(0,0,10,5); i1.DUP_ARR_RGE (1,20,0) END
---	--

DESCRIPTION

L_ARRAY1_RANGE is used in place of BASIC_ARRAY_RANGE, so that a range of arrays may create a set of function type abstract variables when operations are required to perform complete array initialization.

It also allows performing operations that use parts of two different arrays.

The index and range sets are intervals so that it is possible to indicate only portions of these sets without listing all elements involved.

MACHINE PARAMETERS

L_ARRAY1_RANGE (LAUR_minrge, LAUR_maxrge, LAUR_maxidx, LAUR_VALUE):
 The range interval is the LAUR_minrge..LAUR_maxrge interval, the index interval is 0..LAUR_maxidx and LAUR_VALUE is the set of possible values.

VAL_ARR_RGE

syntax $vv \leftarrow \text{VAL_ARR_RGE}(\text{range}, \text{index})$
preconditions range must belong to LAUR_minrge..LAUR_maxrge and index belong to 0..LAUR_maxidx.
outputs vv is a LAUR_VALUE, it is the value of the array range at the index position.

STR_ARR_RGE

syntax $\text{STR_ARR_RGE}(\text{range}, \text{index}, \text{value})$
preconditions range must belong to LAUR_minrge..LAUR_maxrge, index belong to 0..LAUR_maxidx and value belong to LAUR_VALUE.

The value data value is stored in the indexed array range.

COP_ARR_RGE

syntax $\text{COP_ARR_RGE}(\text{dest}, \text{src})$
preconditions dest and src are in LAUR_minrge..LAUR_maxrge

The src array is copied to the dest array.

CMP_ARR_RGE

syntax $bb \leftarrow \text{CMP_ARR_RGE}(\text{range1}, \text{range2})$
preconditions range1 and range2 are in LAUR_minrge..LAUR_maxrge
outputs bb is a BOOL element that is TRUE if the two arrays are equal and FALSE if not.

SET_ARR_RGE

syntax $\text{SET_ARR_RGE}(\text{range}, ii, jj, vv)$
preconditions range must belong to LAUR_minrge..LAUR_maxrge, ii..jj be included in 0..LAUR_maxidx and vv belong to LAUR_VALUE. For implementation reasons, jj must also be different from MAXINT.

The vv value is stored in the array range for all index values between ii and jj. If ii>jj, the array remains unchanged.

DUP_ARR_RGE

syntax $\text{DUP_ARR_RGE}(\text{dest1}, \text{dest2}, \text{src})$
preconditions dest1, dest2, src are in LAUR_minrge..LAUR_maxrge. For implementation reasons, dest2 must also be different from MAXINT.

The src array is duplicated in all of the arrays of the dest1..dest2 interval.

PCOP_ARR_RGE

syntax PCOP_ARR_RGE (dest, idx_dst, src, ii, jj)

preconditions dest and src must be different elements of LAUR_minrge..LAUR_maxrge, ii..jj be a non empty interval of 0..LAUR_maxidx, idx_dst belong to 0..LAUR_maxidx, jj be different from MAXINT and idx_dst + jj - ii belong to 0..LAUR_maxidx (condition necessary to ensure that the copy does not overflow).

The ii..jj part of the src array is copied to the dest array, from the idx_dst index.

PCMP_ARR_RGE

syntax idx, bb ← PCMP_ARR_RGE (rng2, idx2, rng1, ii, jj)

preconditions rng1 and rng2 must belong to LAUR_minrge..LAUR_maxrge, ii..jj be a non empty interval of 0..LAUR_maxidx, idx2 and idx2 + jj-ii are in 0..LAUR_maxidx.

The ii..jj part of array rng1 is compared to the part with the same size in the rng2 array. The $\text{idx2} + \text{jj} - \text{ii} \in 0..LAUR_maxidx$ condition guarantees that this comparison is possible. bb is a Boolean element that is FALSE if the two parts are equal and TRUE if they are different. In the latter case, idx and index are the first element that is different from ii..jj.

IMPORTS REQUIRED

(instances to import as the implementation tree for this library machine sees them with SEES) BASIC_ARITHMETIC, BASIC_BOOL.

WARNING: The implementation of this machine creates the default instance for the BASIC_ARRAY_RANGE machine (IMPORTS BASIC_ARRAY_RANGE(...) clause). Therefore if another instance is necessary, it must be given the name of a non empty instance (for example: i1.BASIC_ARRAY_RANGE).

5.9 L_ARRAY3_RANGE: A Range of Arrays of the Same Size, with Non Ordered Values, Maximum Operations

OPERATIONS

VAL_ARR_RGE	value of an element (promoted operation).
STR_ARR_RGE	write an element (promoted operation).
COP_ARR_RGE	copy an array to another (promoted operation).
CMP_ARR_RGE	compare two arrays (promoted operation).
DUP_ARR_RGE	duplicate the same array to a set of arrays (promoted operation).
SET_ARR_RGE	copy the same value to an index interval in one of the arrays (promoted operation).
PCOP_ARR_RGE	copy part of one of the arrays to a different array, at a given position (promoted operation).
PCMP_ARR_RGE	search for the first element that is different between two parts of two arrays. A Boolean element indicates whether this element was found and, in this case, the index of this element is returned (promoted operation).
SWAP_RGE	exchange two array elements.
RIGHT_SHIFT_RGE	shift part of an array to the large index.
LEFT_SHIFT_RGE	shift part of an array to the small index.
SEARCH_MAX_EQL_RGE	search for the last element that equals a value in part of an array.
SEARCH_MIN_EQL_RGE	search for the first element that equals a value in part of an array.
REVERSE_RGE	reverse the order of the elements in part of an array.

EXAMPLE

The following example is a machine that represents the color assigned to 101 dots for each array in a range; this color may be red, green or blue for each dot. A operation enables finding a red dot in an array.

MACHINE m3 SETS COLOR = {red,green,blue} VARIABLES color INVARIANT color $\in 0..10 \rightarrow (0..100 \rightarrow \text{COLOR})$ INITIALISATION color := (0..10) \times {(0..100) \times {red}} OPERATIONS ii,bb \leftarrow find_red(rng) = PRE rng $\in 0..10 \wedge$ rouge $\in \text{ran}(\text{color}(\text{rng}))$ THEN ii: $\in \text{color}(\text{rng})^{-1}[\{\text{red}\}] \parallel$ bb: $\in \text{BOOL}$ END END	IMPLEMENTATION m3_1 REFINES m3 IMPORTS i1.L_ARRAY3_RANGE(0,10,100,COLOR) INVARIANT i1.arr_rge = color INITIALISATION i1.SET_ARR_RGE(0,0,100,red); i1.DUP_ARR_RGE(1,10,0) OPERATIONS ii,bb \leftarrow find_red(rng) = VAR bb IN ii,bb \leftarrow i1.SEARCH_MAX_EQL_RGE(rng,0,100,red) END END
---	--

DESCRIPTION

L_ARRAY3_RANGE is the most complete of the two dimensional array machines with no constraint³. This makes it possible to create arrays with values that are the elements of an enumerated set, while retaining access to complete operations such as reversing the order of elements.

The operation that is not available is the one that would require an order relation on the elements in the array: sort.

MACHINE PARAMETERS

L_ARRAY3_RANGE (LATR_minrge, LATR_maxrge, LATR_maxidx, LATR_VALUE):

The range interval is LATR_minrge..LATR_maxrge, the index interval 0..LATR_maxidx and LATR_VALUE is the set of possible values.

VAL_ARR_RGE

syntax vv \leftarrow VAL_ARR_RGE (range, index)

preconditions range must belong to LATR_minrge..LATR_maxrge, index belong to 0..LATR_maxidx

outputs vv is a LATR_VALUE, it is the value of the array range at the index position.

STR_ARR_RGE

syntax STR_ARR_RGE (range, index, value)

preconditions range must belong to LATR_minrge..LATR_maxrge, index belong to 0..LATR_maxidx and value belong to LATR_VALUE.

The LATR_VALUE value is stored in the array range in the index.

³L_ARRAY5_RANGE can only have a finite integer set as range.

COP_ARR_RGE

syntax COP_ARR_RGE (dest, src)

preconditions dest and src are in LATR_minrge..LATR_maxrge

The src array is copied to the dest array.

CMP_ARR_RGE

syntax bb \leftarrow CMP_ARR_RGE (range1, range2)

preconditions range1 and range2 are in LATR_minrge..LATR_maxrge

outputs bb is an BOOL that equals TRUE if the two arrays are equal and FALSE if not.

SET_ARR_RGE

syntax SET_ARR_RGE (range, ii, jj, vv)

preconditions range must belong to LATR_minrge..LATR_maxrge, ii..jj be included in 0..LATR_maxidx and vv belong to LATR_VALUE. For implementation reasons, jj must also be different to MAXINT.

Value vv is stored in the array range for all indexes in the range from ii to jj. If ii > jj, the array remains unchanged.

DUP_ARR_RGE

syntax DUP_ARR_RGE (dest1, dest2, src)

preconditions dest1, dest2, src are in LATR_minrge..LATR_maxrge. For implementation reasons, dest2 must also be different to MAXINT.

The src array is duplicated in all of the arrays of the dest1..dest2 interval.

PCOP_ARR_RGE

syntax PCOP_ARR_RGE (dest, idx_dst, src, ii, jj)

preconditions dest and src must belong to LATR_minrge..LATR_maxrge and be different, ii..jj be a non empty interval of 0..LATR_maxidx, idx_dst belong to 0..LATR_maxidx, jj be different from MAXINT and idx_dst + jj - ii belong to 0..LATR_maxidx (necessary condition to avoid copy overflow).

The ii..jj part in the src array is copied to the dest array, from the idx_dst index.

PCMP_ARR_RGE

syntax idx, bb \leftarrow PCMP_ARR_RGE (rng2, idx2, rng1, ii, jj)

preconditions rng1 and rng2 are in LATR_minrge..LATR_maxrge, ii..jj is a non empty interval of 0..LATR_maxidx idx2 and idx2 + jj - ii are in 0..LATR_maxidx.

The ii..jj part of array rng1 is compared with the part with the same size in array rng2. The idx2 + jj - ii \in 0..LATR_maxidx condition guarantees that this comparison is possible. bb is a Boolean element that is FALSE if the two parts are equal and TRUE if they are different. In the latter case, idx is the index of the first element that is different to ii..jj.

SWAP_RGE

syntax SWAP_RGE (rng,ii,jj)

preconditions rng is in LATR_minrge..LATR_maxrge, ii and jj in 0..LATR_maxidx.

The ii and jj elements in the array are exchanged.

RIGHT_SHIFT_RGE

syntax RIGHT_SHIFT_RGE (rng,ii,jj,nn)

preconditions rng must belong to LATR_minrge..LATR_maxrge, ii, jj and nn belong to 0..LATR_maxidx, with $ii \leq jj$ and $jj+nn \leq \text{LATR_maxidx}$ to allow a right shift by nn spaces.

The ii+nn..jj+nn part in the rng array receives a copy of the ii..jj part of this same array (shift right by nn spaces).

LEFT_SHIFT_RGE

syntax LEFT_SHIFT_RGE (rng,ii,jj,nn)

preconditions rng is in LATR_minrge..LATR_maxrge; ii,jj must be in 0..LATR_maxidx, with $ii \leq jj$. nn must be a NAT with $nn \leq ii$ to allow the left shift by nn spaces. For implementation reasons, jj must be equal to MAXINT.

The ii-nn..jj-nn part of the rng array receives a copy of the ii..jj part of this same array (shift left by nn spaces).

SEARCH_MAX_EQL_RGE

syntax rr,bb \leftarrow SEARCH_MAX_EQL_RGE (rng,ii,jj,vv)

preconditions rng must be in LATR_minrge..LATR_maxrge. ii and jj must be in 0..LATR_maxidx, $ii \leq jj$ and vv must belong to LATR_VALUE.

outputs TRUE if vv was found, FALSE if not. rr is a NAT, if bb = TRUE then rr is the largest index in the rng array equal to vv.

Search for an element in an array equal to vv, by scanning the ii..jj part starting from jj.

SEARCH_MIN_EQL_RGE

syntax rr,bb \leftarrow SEARCH_MIN_EQL_RGE (rng,ii,jj,vv)

preconditions rng must belong to LATR_minrge..LATR_maxrge, ii and jj belong to 0..LATR_maxidx, $ii \leq jj$ and vv belong to LATR_VALUE.

outputs TRUE if vv was found, FALSE if not. rr is a NAT, if bb = TRUE, then rr is the smallest index in the rng array equal to vv.

Search for an element in an array that is equal to vv, by scanning the ii..jj part starting from ii.

REVERSE_RGE

syntax REVERSE_RGE(rng,ii,jj)

preconditions rng must belong to LATR_minrge..LATR_maxrge, ii and jj belong to 0..LATR_maxidx.

Reversing the order of elements in the ii..jj part of the rng array.

IMPORTS REQUIRED

(instances to import as the implementation tree for this library machine sees them with SEES) BASIC_ARITHMETIC; BASIC_BOOL.

WARNING: The implementation of this machine creates the default instance for the BASIC_ARRAY_RANGE machine (IMPORTS BASIC_ARRAY_RANGE(...) clause). Therefore if another instance is necessary, it must be given a non empty instance name (for example: i1.BASIC_ARRAY_RANGE)

5.10 L_ARRAY5_RANGE: Range of Arrays of the Same Size, with Ordered Value Numerical Indexes, Sort Operation

OPERATIONS

VAL_ARR_RGE	value of an element (promoted operation).
STR_ARR_RGE	write an element (promoted operation).
COP_ARR_RGE	copy an array to another (promoted operation).
CMP_ARR_RGE	compare two arrays (promoted operation).
DUP_ARR_RGE	duplicate the same array to a set of arrays (promoted operation).
SET_ARR_RGE	copy the same value to an index range in one of the arrayx (promoted operation).
PCOP_ARR_RGE	copy part of one of the arrays to a different array, to a given position (promoted operation).
PCMP_ARR_RGE	search for the first different element between two parts of two arrays. A Boolean element indicates whether this element was found and, in this case, the index of this element is returned (promoted operation).
SWAP_RGE	exchange two elements in an array (promoted operation).
RIGHT_SHIFT_RGE	shift a part of an array to the large index (promoted operation).
LEFT_SHIFT_RGE	shift a part of an array to the small index (promoted operation).
SEARCH_MAX_EQL_RGE	search for the last element that is equal to a value in an array range (promoted operation).
SEARCH_MIN_EQL_RGE	search for the first element that equals a value in an array range (promoted operation).
REVERSE_RGE	reverse the order of the elements of a part of an array (promoted operation).
SEARCH_MIN_GEQ_RGE	search for the first element that exceeds a value in an array range.
ASCENDING_SORT_RGE	sort part of an array and arrange in ascending order.

EXAMPLE

MACHINE m5 VARIABLES vv INVARIANT $vv \in 0..20 \rightarrow (0..4 \rightarrow 0..255) \wedge$ $\forall (xx,yy).(xx \in 0..20 \wedge yy \in 0..3$ $\Rightarrow vv(yy)(xx) \geq vv(yy)(xx+1))$ INITIALISATION $vv : (vv \in 0..20 \rightarrow (0..4 \rightarrow 0..255) \wedge$ $\forall (xx,yy).(xx \in 0..20 \wedge yy \in 0..3 \Rightarrow$ $vv(yy)(xx) \geq vv(yy)(xx+1)))$ END	IMPLEMENTATION m5_1 REFINES m5 IMPORTS L_ARRAY5_RANGE(0,20,4,0,255) INVARIANT arr_rge = vv INITIALISATION SET_ARR_RGE(0,0,4,50); STR_ARR_RGE(0,2,10); STR_ARR_RGE(0,4,30); ASCENDING_SORT_RGE(0,0,4); REVERSE_RGE(0,0,4); DUP_ARR_RGE(1,20,0) END
--	--

DESCRIPTION

L_ARRAY5_RANGE is the most complete two dimensional array machines. It especially contains a sort operation, implanted by a shift sort (fast algorithm).

MACHINE PARAMETERS

L_ARRAY5_RANGE (LACR_minrge, LACR_maxrge, LACR_maxidx, LACR_minval, LACR_maxval):

LACR_minrge..LACR_maxrge is the set of ranges, 0..LACR_maxidx is the set of indexes and LACR_minval..LACR_maxval, the set of possible values. All of the parameters must be NATs: this machine does not allow negative values.

In addition, $LACR_minrge \leq LACR_maxrge$, $1 \leq LACR_maxidx$ and $LACR_minval \leq LACR_maxval$.

VAL_ARR_RGE

syntax $vv \leftarrow VAL_ARR_RGE \text{ (range, index)}$

preconditions range must belong to LACR_minrge..LACR_maxrge, index belong to 0..LACR_maxidx.

outputs vv is a LACR_VALUE, it is the value of the array range at the index position.

STR_ARR_RGE

syntax $STR_ARR_RGE \text{ (range, index, value)}$

preconditions range must be in LACR_minrge..LACR_maxrge index must be in 0 .. LACR_maxidx value must belong to LACR_VALUE.

The value of the value element is stored in the array range as an index.

COP_ARR_RGE

syntax COP_ARR_RGE (dest, src)
preconditions dest and src are in LACR_minrge..LACR_maxrge

The src array is copied to the dest array.

CMP_ARR_RGE

syntax $bb \leftarrow \text{CMP_ARR_RGE}(\text{range1}, \text{range2})$
preconditions range1 and range2 are in LACR_minrge..LACR_maxrge
outputs bb is a BOOL element that is TRUE if the two arrays are equal and FALSE if not.

SET_ARR_RGE

syntax SET_ARR_RGE (range, ii, jj, vv)
preconditions range must belong to LACR_minrge..LACR_maxrge, ii..jj be included in 0..LACR_maxidx and vv belong to LACR_VALUE. For implementation reasons, it is also necessary that jj be different from MAXINT.

The vv value is stored in the array range for all indexes between ii and jj. If ii>jj, the array remains unchanged.

DUP_ARR_RGE

syntax DUP_ARR_RGE (dest1, dest2, src)
preconditions dest1, dest2, src are in LACR_minrge..LACR_maxrge. For implementation reasons, it is also necessary for dest2 to be different from MAXINT.

The src array is duplicated to all arrays for the dest1..dest2 range.

PCOP_ARR_RGE

syntax PCOP_ARR_RGE (dest, idx_dst, src, ii, jj)
preconditions dest and src must be different elements of LACR_minrge..LACR_maxrge, ii..jj be a non empty subset of 0..LACR_maxidx and idx_dst belong to 0..LACR_maxidx; jj is different from MAXINT and idx_dst + jj - ii belong to 0..LACR_maxidx (condition to avoid copy overflow).

The ii..jj range in the src array is copied to the dest array, for the idx_dst index.

PCMP_ARR_RGE

syntax $\text{idx}, bb \leftarrow \text{PCMP_ARR_RGE}(\text{rng2}, \text{idx2}, \text{rng1}, \text{ii}, \text{jj})$
preconditions rng1 and rng2 are in LACR_minrge..LACR_maxrge, ii..jj is a non empty range 0..LACR_maxidx, idx2 and idx2 + jj-ii are in 0..LACR_maxidx.

The ii..jj part of the rng1 array is compared with the part of the same size in the rng2 array. The $\text{idx2} + \text{jj} - \text{ii} \in 0..LACR_maxidx$ condition guarantees that this comparison is possible. bb is a Boolean element that is FALSE if the two parts are equal and TRUE if they are different. In the latter case, idx is the index of the first element that is different from ii..jj.

SWAP_RGE

syntax SWAP_RGE (rng,ii,jj)

preconditions rng is in LACR_minrge..LACR_maxrge, ii and jj in 0..LACR_maxidx.

The ii and jj elements in the array are exchanged.

RIGHT_SHIFT_RGE

syntax RIGHT_SHIFT_RGE (rng,ii,jj,nn)

preconditions rng must belong to LACR_minrge..LACR_maxrge. ii, jj and nn belong to 0..LACR_maxidx, with $ii \leq jj$ and $jj+nn \leq \text{LACR_maxidx}$ to allow the shift right by nn spaces.

The ii+nn..jj+nn part of the rng array receives a copy of the ii..jj part from this same array (shift nn spaces to the right).

LEFT_SHIFT_RGE

syntax LEFT_SHIFT_RGE (rng,ii,jj,nn)

preconditions rng must belong to LACR_minrge..LACR_maxrge, ii and jj belong to 0..LACR_maxidx, with $ii \leq jj$. nn must belong to NAT with $nn \leq ii$ to make possible the left shift by nn spaces. For implementation reasons, jj cannot equal MAXINT.

The ii-nn..jj-nn part of the rng array receives a copy of the ii..jj part of this same array (shift nn spaces to the left).

SEARCH_MAX_EQL_RGE

syntax rr,bb \leftarrow SEARCH_MAX_EQL_RGE (rng,ii,jj,vv)

preconditions rng must belong to LACR_minrge..LACR_maxrge. ii and jj belong to 0..LACR_maxidx, $ii \leq jj$ and vv must belong to LACR_VALUE.

outputs TRUE if vv was found, FALSE if not. rr is a NAT, if bb = TRUE then rr the largest index in the array that equals vv.

Search for an array element that equals vv, by scanning the ii..jj part starting from jj.

SEARCH_MIN_EQL_RGE

syntax rr,bb \leftarrow SEARCH_MIN_EQL_RGE (rng,ii,jj,vv)

preconditions rng must belong to LACR_minrge..LACR_maxrge, ii and jj belong to 0..LACR_maxidx, $ii \leq jj$ and vv must belong to VALUE.

outputs TRUE if vv was found, FALSE if not. rr is a NAT, if bb = TRUE, then rr is the smallest index in the rng array equal to vv.

Search for an element in an array equal to vv, by scanning the ii..jj part starting from ii.

REVERSE_RGE

syntax REVERSE_RGE(rng,ii,jj)

preconditions rng must belong to LACR_minrge..LACR_maxrge, ii and jj belong to 0..LACR_maxidx.

Reverse the order of elements in the ii..jj range of the rng array.

SEARCH_MIN_GEQ_RGE

syntax $ii, bb \leftarrow \text{SEARCH_MIN_GEQ_RGE}(rng, jj, kk, vv)$

preconditions rng must belong to $LACR_minrge..LACR_maxrge$. jj and kk belong to $0..LACR_maxidx$, $jj \leq kk$ and vv belong to $LACR_minval..LACR_maxval$. For implementation reasons, kk must be different from $MAXINT$.

outputs bb is a Boolean element, $TRUE$ is an element that exceeds or is equal to the vv value found, $FALSE$ if not. ii is a NAT , if $bb = TRUE$, then ii is the smallest index in the image array that exceeds or is equal to vv .

Search for an element that exceeds or is equal to vv in the $jj..kk$ range, starting from jj .

ASCENDING_SORT_RGE

syntax $\text{ASCENDING_SORT_RGE}(rng, ii, jj)$

preconditions rng must belong to $LACR_minrge..LACR_maxrge$, ii and jj belong to $0..LACR_maxidx$. For implementation reasons, ii and jj must not be different from $MAXINT$.

Shift sort, in ascending order (starting with the smallest) on the $ii..jj$ range in an array.

IMPORTS REQUIRED

(instances to import as the implementation tree for this library machine sees them with SEES)

BASIC_ARITHMETIC; BASIC_BOOL.

WARNING: The implementation of this machine creates the default instance for the **BASIC_ARRAY_RANGE** machine (**IMPORTS BASIC_ARRAY_RANGE(...)** clause). Therefore if another instance is necessary, it must be given a non empty instance name (for example: $i1.BASIC_ARRAY_RANGE$)

5.11 L_SEQUENCE_RANGE: Range of Sequences

OPERATIONS

LEN_SEQ_RGE	gives the size of a sequence.
IS_FULL_SEQ_RGE	indicates whether a sequence is full.
IS_INDEX_SEQ_RGE	indicates whether an integer is in the sequence domain.
VAL_SEQ_RGE	gives the value of a sequence for a valid index.
FIRST_SEQ_RGE	gives the first element of a sequence.
LAST_SEQ_RGE	gives the last element of a sequence.
PUSH_SEQ_RGE	adds an element to a sequence.
POP_SEQ_RGE	removes the last element from a sequence.
STR_SEQ_RGE	changes the value of an element in a sequence.
RMV_SEQ_RGE	removes an element from a sequence, the size of which is reduced by 1.
INS_SEQ_RGE	adds an element to a sequence, the size of which increases by 1.
CLR_SEQ_RGE	empties a sequence.
TAIL_SEQ_RGE	removes the first element from a sequence.
KEEP_SEQ_RGE	only retains the first N in a sequence elements.
CUT_SEQ_RGE	cuts the N first elements from a sequence.
PART_SEQ_RGE	only retains in a sequence the indexes between the two limit values.
REV_SEQ_RGE	reverses the order of the elements in a sequence.
FIND_FIRST_SEQ_RGE	searches for a value in a sequence, returns a Boolean element indicating whether it was found, and if yes, returns the smallest corresponding index.
FIND_LAST_SEQ_RGE	searches for a value in a sequence, returns a Boolean element indicating whether it was found and if yes, returns the largest corresponding index.
COP_SEQ_RGE	copies one of the sequences to another.
CMP_SEQ_RGE	compares two sequences.
PCOP_SEQ_RGE	partial copy from one sequence to another.
PCMP_SEQ_RGE	partial comparison of two sequences.

EXAMPLE

The example below shows the use of `L_SEQUENCE_RANGE` on a numbered set.

<pre> MACHINE sr SETS ST = {classical,baroque,rock,rap,funk} VARIABLES vv INVARIANT vv ∈ 1..5 → seq(ST) ∧ ∀rr.(rr ∈ 1..5 ⇒ size(vv(rr)) ≤ 10) INITIALISATION vv:=(1..5) × {[baroque,rock,rap]} OPERATIONS ii,bb ← trouve_rap(rng) = PRE rng ∈ 1..5 THEN ii:∈ vv(rng)⁻¹[[rap]] bb:∈ BOOL END END </pre>	<pre> IMPLEMENTATION sr_l REFINES sr IMPORTS s1.L_SEQUENCE_RANGE(1,5,10,ST) INVARIANT s1.seq_rge = vv INITIALISATION s1.CLR_SEQ_RGE(1); s1.PUSH_SEQ_RGE(1,baroque); s1.PUSH_SEQ_RGE(1,rock); s1.PUSH_SEQ_RGE(1,rap); s1.COP_SEQ_RGE(2,1); s1.COP_SEQ_RGE(3,1); s1.COP_SEQ_RGE(4,1); s1.COP_SEQ_RGE(5,1) OPERATIONS ii,bb ← trouve_rap(rng) = BEGIN bb,ii ← s1.FIND_FIRST_SEQ_RGE(rng,rap) END END </pre>
--	---

DESCRIPTION

`L_SEQUENCE_RANGE` enables implementing and using a set number of sequences with a fixed maximum size. The sequence number evolves in a range that is a machine parameter, the maximum size of all of the sequences is also a machine parameter. The purpose is to be able to make comparisons and copies between these sequences directly, using an additional operation to the traditional operations on each of the sequences.

MACHINE PARAMETERS

`L_SEQUENCE_RANGE` (`LSR_minrge`, `LSR_maxrge`, `LSR_maxsize`, `LSR_VALUE`): the variable is a total function of `LSR_minrge..LSR_maxrge` in the set of `VALUE` sequences with a maximum size of `LSR_maxsize`.

LEN_SEQ_RGE

syntax `nn ← LEN_SEQ_RGE (range)`

preconditions range must belong to the `LSR_minrge..LSR_maxrge` range.

outputs nn is the size of the range position , `nn ∈ 0..LSR_maxsize`.

Gives the size of a sequence.

IS_FULL_SEQ_RGE

- syntax* $bb \leftarrow \text{IS_FULL_SEQ_RGE}(\text{range})$
- preconditions* range must belong to the range `LSR_minrge..LSR_maxrge`.
- outputs* bb is TRUE if the range position sequence is full, FALSE if not.

Indicates whether a sequence is full.

IS_INDEX_SEQ_RGE

- syntax* $bb \leftarrow \text{IS_INDEX_SEQ_RGE}(\text{range}, ii)$
- preconditions* range must belong to the `LSR_minrge..LSR_maxrge` range, ii must be a NAT.
- outputs* bb is TRUE if ii is an index in the range position sequence, FALSE if not.

Identifies whether an integer is in a sequence domain.

VAL_SEQ_RGE

- syntax* $vv \leftarrow \text{VAL_SEQ_RGE}(\text{range}, ii)$
- preconditions* range must belong to the `LSR_minrge..LSR_maxrge` range, ii must be an index in the range position sequence ($ii \in 1..size(\text{seq_rge}(\text{range}))$).
- outputs* vv is the value of the ii-th element in the range position sequence ($vv \in \text{VALUE}$).

Gives the value of a sequence for a valid index.

FIRST_SEQ_RGE

- syntax* $vv \leftarrow \text{FIRST_SEQ_RGE}(\text{range})$
- preconditions* range must belong to the `LSR_minrge..LSR_maxrge` range, the range position sequence must not be empty.
- outputs* vv is the value of the first element in the range position sequence ($vv \in \text{VALUE}$).

Gives the first element in a sequence.

LAST_SEQ_RGE

- syntax* $vv \leftarrow \text{LAST_SEQ_RGE}(\text{range})$
- preconditions* range must be in the `LSR_minrge..LSR_maxrge` range, the range position sequence must not be empty.
- outputs* vv is the value of the last element in the range position sequence ($vv \in \text{VALUE}$).

Gives the last element of a sequence.

PUSH_SEQ_RGE

- syntax* $\text{PUSH_SEQ_RGE}(\text{range}, vv)$
- preconditions* range must belong to the `LSR_minrge..LSR_maxrge` range, vv must be in `LSR_VALUE` and the range position sequence cannot be full.

Adds an element to a sequence.

POP_SEQ_RGE

syntax POP_SEQ_RGE (range)

preconditions range must belong to the LSR_minrge..LSR_maxrge range, the range position sequence must not be empty.

Removes the last element in a sequence.

STR_SEQ_RGE

syntax STR_SEQ_RGE (range, ii, vv)

preconditions range must belong to LSR_minrge..LSR_maxrge, ii be a valid index in the range position sequence and vv belong to LSR_VALUE.

Change the value of an element in a sequence.

RMV_SEQ_RGE

syntax RMV_SEQ_RGE (range, ii)

preconditions range must belong to the LSR_minrge..LSR_maxrge range, ii must be a valid index in the range sequence.

Removes an element from a sequence, the size of which decreases by 1.

INS_AFT_SEQ_RGE

syntax INS_AFT_SEQ_RGE (range, ii, vv)

preconditions range must belong to the LSR_minrge..LSR_maxrge range, ii must be a valid index in the range position sequence, vv must be in LSR_VALUE, the range position sequence must not be full.

Adds an element to a sequence, the size of which increases by 1.

CLR_SEQ_RANGE

syntax CLR_SEQ_RANGE (range)

preconditions range must belong to the LSR_minrge..LSR_maxrge range.

Clears a sequence.

TAIL_SEQ_RGE

syntax TAIL_SEQ_RGE (range)

preconditions range must belong to the LSR_minrge..LSR_maxrge range and the range position sequence cannot be empty.

Removes the first element in a sequence.

KEEP_SEQ_RGE

syntax KEEP_SEQ_RGE (range, nn)

preconditions range must belong to the LSR_minrge..LSR_maxrge range, nn must be a NAT.

Only retains the nn first elements in a sequence. For $nn = \text{size}(\text{seq_rge}(\text{range}))$; this operation has no effect.

CUT_SEQ_RGE

syntax CUT_SEQ_RGE (range, nn)

preconditions range must belong to the LSR_minrge..LSR_maxrge range, nn must be in NAT.

Clears the sequence of its first nn elements. For $nn = \text{size}(\text{seq_rge}(\text{range}))$, this operation is equivalent to CLR_SEQ_RGE.

PART_SEQ_RGE

syntax PART_SEQ_RGE (range, ii, jj)

preconditions range must belong to the LSR_minrge..LSR_maxrge range, ii and jj must be NATs that are not null, with $ii \leq jj$.

In a sequence, only retains the indexes between two limits. ii..jj may not be in the sequence domain.

REV_SEQ_RGE

syntax REV_SEQ_RGE (range)

preconditions range must belong to the LSR_minrge..LSR_maxrge range.

Reverses the order of the elements in a sequence.

FIND_FIRST_SEQ_RGE

syntax bb, ii \leftarrow FIND_FIRST_SEQ_RGE (range, vv)

preconditions range must belong to the LSR_minrge..LSR_maxrge range, vv must be in LSR_VALUE.

outputs bb is TRUE if vv is in the range position sequence, FALSE if not. ii is a NAT, if bb = TRUE, it indicates the first position that equals vv in the sequence.

Searches for a value in a sequence starting from the beginning.

FIND_LAST_SEQ_RGE

syntax bb, ii \leftarrow FIND_LAST_SEQ_RGE (range, vv)

preconditions range must belong to the LSR_minrge..LSR_maxrge range, vv must be in LSR_VALUE.

outputs bb is TRUE if vv is in the range position sequence, FALSE if not. ii is a NAT; if bb = TRUE, this indicates the last position that equals vv in the sequence.

Searches for a value in a sequence, starting from the end.

COP_SEQ_RGE

syntax COP_SEQ_RGE (dst, src)

preconditions dst and src must belong to the LSR_minrge..LSR_maxrge range.

Copy the seq_rge(src) sequence to the seq_rge(dst) sequence.

CMP_SEQ_RGE

syntax $bb \leftarrow \text{CMP_SEQ_RGE}(\text{rng1}, \text{rng2})$

preconditions rng1 and rng2 must belong to the $\text{LSR_minrge}..\text{LSR_maxrge}$ range.

outputs bb is TRUE if the two rng1 and rng2 position sequences are equal, FALSE if not.

Compare two sequences.

PCOP_SEQ_RGE

syntax $\text{PCOP_SEQ_RGE}(\text{dst}, \text{idx}, \text{src}, \text{ii}, \text{jj})$

preconditions dst and src must belong to the $\text{LSR_minrge}..\text{LSR_maxrge}$ range, dst must be different from src , ii and jj must be valid indexes in the src position sequence, with $\text{ii} \leq \text{jj}$ and $\text{jj} \leq \text{MAXINT}-1$ idx must be a valid index for the dst sequence or where the size of this sequence $+1$, $\text{idx} + \text{jj} - \text{ii}$ belongs to the $1..\text{LSR_maxsize}$ range.

Copy the $\text{ii}..\text{jj}$ part of the src position sequence to the dst position from the idx index.

PCMP_SEQ_RGE

syntax $\text{idx}, bb \leftarrow \text{PCMP_SEQ_RGE}(\text{rng1}, \text{ii}, \text{jj}, \text{rng2}, \text{kk})$

preconditions rng1 and rng2 must be in the $\text{LSR_minrge}..\text{LSR_maxrge}$ range, ii and jj must be valid indexes in the rng1 and $\text{ii} \leq \text{jj}$ position sequences, kk must be a valid index in the rng2 position sequence, $(\text{kk} + \text{jj} - \text{ii})$ must be a valid index in the rng2 position sequence.

output bb is TRUE if there is an element of the $\text{ii}..\text{jj}$ part in the $\text{seq_rge}(\text{rng1})$ sequence that is different to the $\text{kk}..(\text{kk} + \text{jj} - \text{ii})$ part of the $\text{seq_rge}(\text{rng2})$ sequence, FALSE if not. idx is a NAT if bb is TRUE, the idx represents the index of the first element that is different in the $\text{seq_rge}(\text{rng1} \in \text{ii}..\text{jj})$ sequence.

Partial comparison of two sequences.

IMPORTS REQUIRED

(instances to import as the implementation tree for this library machine sees them with SEES)

BASIC_ARITHMETIC ;

BASIC_BOOL.

WARNING: The implementation of this machine creates the default instance for the BASIC_ARRAY_RANGE and BASIC_ARRAY_VAR machines. Therefore, if other instances are required they must be given a name that is not blank.

(example: $\text{i1.BASIC_ARRAY_RANGE}$).

5.12 L_ARRAY_COLLECTION: collection of arrays of the same size

OPERATIONS

CRE_ARR_COL	returns a Boolean element that indicates that there remains an array available in the collection and gives the index of this available array.
DEL_ARR_COL	releases the specified array.
VAL_ARR_COL	read an element from one of the valid arrays.
STR_ARR_COL	write an element from one of the valid arrays.
COP_ARR_COL	copy one of the arrays to another.
CMP_ARR_COL	compare two arrays.

EXAMPLE

<p>MACHINE m1 OPERATIONS ii1,ii2 \leftarrow initialise_arrayx(vv) = PRE vv \in 1..10 THEN ii1:\in NAT ii2:\in NAT END END</p>	<p>IMPLEMENTATION M1_1 REFINES M1 IMPORTS L_ARRAY_COLLECTION(4,1..10,1..10) OPERATIONS ii1,ii2 \leftarrow initialise_arrayx(vv) = BEGIN VAR b1,b2 IN ii1,b1 \leftarrow CRE_ARR_COL; ii2,b2 \leftarrow CRE_ARR_COL; STR_ARR_COL(ii1,1,vv); COP_ARR_COL(ii2,ii1) END END END</p>
--	---

DESCRIPTION

L_ARRAY_COLLECTION is used to handle identically sized one dimensional arrays. It contains basic operations (create, delete, read, write, compare).

MACHINE PARAMETERS

L_ARRAY_COLLECTION (LACOLL_maxobj, LACOLL_INDEX, LACOLL_VALUE):
LACOLL_maxobj is the maximum number of arrays in the collection. LACOLL_INDEX is the set of array indexes, LACOLL_VALUE is the set of possible values of array elements.

CRE_ARR_COL

Syntax $ii, bb \leftarrow \text{CRE_ARR_COL}$

Outputs bb is a Boolean element indicating whether any available arrays are left in the collection, ii is the index of this available array.

Assigning an array in the collection.

DER_ARR_COL

Syntax $\text{DEL_ARR_COL } (ii)$

Preconditions ii must belong to $1..\text{LACOLL_maxobj}$

The array of index ii in the collection is released. It may once again be assigned using CRE_ARR_COL .

VAL_ARR_COL

Syntax $vv \leftarrow \text{VAL_ARR_COL } (ii, jj)$

Preconditions ii must belong to $1..\text{LACOLL_maxobj}$ and jj belong to LACOLL_INDEX .

Output vv contains the jj number value of array ii .

Use vv to store the value of number jj in array ii .

STR_ARR_COL

Syntax $\text{STR_ARR_COL } (ii, jj, vv)$

Preconditions ii must belong to $1..\text{LACOLL_maxobj}$, jj belong to LACOLL_INDEX and vv belong to LACOLL_VALUE .

Write the value of vv to cell number jj in array ii .

COP_ARR_COL

Syntax $\text{COP_ARR_COL } (\text{dest}, \text{src})$

Preconditions dest and src must belong to $1..\text{LACOLL_maxobj}$.

Copy the contents of the src array to the dest array.

CMP_ARR_COL

Syntax $bb \leftarrow \text{CMP_ARR_COL } (\text{range } 1, \text{range } 2)$

Preconditions $\text{range } 1$ and $\text{range } 2$ must belong to $1..\text{LACOLL_maxobj}$.

Output bb is a Boolean element indicating whether array $\text{range } 1$ and $\text{range } 2$ are identical.

Comparison between the two 2 arrays.

IMPORTS REQUIRED

(instances to import as the implementation tree for this library machine sees them with SEES).

BASIC_ARITHMETIC BASIC_BOOL

5.13 L_ARRAY1_COLLECTION: array of the same size, with numerical indexes

OPERATIONS

CRE_ARR_COL	returns a Boolean element indicating whether an array remains available in the collection and the index of this available array (promoted operation).
DEL_ARR_COL	releases the array mentioned (promoted operation).
VAL_ARR_COL	read an element from one of the valid arrays (promoted operation).
STR_ARR_COL	write an element from one of the valid arrays (promoted operation).
COP_ARR_COL	copy one of the arrays to another (promoted operation).
CMP_ARR_COL	compare two arrays (promoted operation).
SET_ARR_COL	copy the same value to an index range in one of the arrays.
PCOP_ARR_COL	copy part of one of the arrays to another in a given position.
PCMP_ARR_COL	search for the first different element between two parts of two different arrays. A Boolean element indicates whether the element was found and in this case, the index of this element is returned.

EXAMPLE

Using SET_ARR_COL to fill-in two arrays and PCOP_ARR_COL to define a third one. Note the need to test the Boolean output elements from CRE_ARR_COL in order to use the arrays created.

The example is as follows:

```
MACHINE
M1
OPERATIONS
  op = skip
END
```

```

IMPLEMENTATION
  M1_1
REFINES
  M1
IMPORTS
  L_ARRAY1_COLLECTION(3,3,1,10)
OPERATIONS
  op = VAR i1,i2,i3,b1,b2,b3 IN
    i1,b1 ← CRE_ARR_COL;
    i2,b2 ← CRE_ARR_COL;
    i3,b3 ← CRE_ARR_COL;
    IF b1 = TRUE ∧
      b2 = TRUE ∧
      b3 = TRUE
    THEN
      SET_ARR_COL(i1,0,3,1);
      SET_ARR_COL(i2,0,3,2);
      PCOP_ARR_COL(i3,0,i1,0,1);
      PCOP_ARR_COL(i3,2,i2,2,3)
    END
  END
END

```

DESCRIPTION

L_ARRAY1_COLLECTION enables the use of a collection of arrays without the need to code loops to position a set of elements or arrays. This was not possible with the previous machine L_ARRAY_COLLECTION where index sets are normally unordered.

MACHINE PARAMETERS

L_ARRAY1_COLLECTION (LAUC_maxobj, LAUC_maxidx, LAUC_minval, LAUC_maxval): The variable is a partial function of 1..LAUC_maxobj in the set of total functions of 0..LAUC_maxidx to LAUC_minval..LAUC_maxval. LAUC_maxobj is a NAT1 that is different from MAXINT. LAUC_maxidx, LAUC_minval and LAUC_maxval are NATs and $\text{LAUC_minval} \leq \text{LAUC_maxval}$.

CRE_ARR_COL

Syntax ii, bb ← CRE_ARR_COL

Outputs bb is a Boolean element indicating whether any available arrays remain in the collection, ii is the index of this available array.

Allocate an array in the collection.

DEL_ARR_COL

Syntax DEL_ARR_COL (ii)

Preconditions ii must belong to 1..LAUC_maxobj

Array ii in the collection is released. It may once again be assigned using CRE_ARR_COL.

VAL_ARR_COL

Syntax $vv \leftarrow \text{VAL_ARR_COL} (ii, jj)$

Preconditions ii must belong to $1..LAUC_maxobj$ jj must belong to $1..LAUC_maxidx$.

Output vv contains the value of number jj in array ii .

Store in vv the value of number jj in array ii .

STR_ARR_COL

Syntax $\text{STR_ARR_COL} (ii, jj, vv)$

Preconditions ii must belong to $1..LAUC_maxobj$; jj must belong to $1..LAUC_maxidx$.
 vv must belong to $LAUC_VALUE$.

Write value vv to the jj th cell in array ii .

COP_ARR_COL

Syntax $\text{COP_ARR_COL} (\text{dest}, \text{src})$

Preconditions dest and src must belong to $1..LAUC_maxobj$.

Copy the contents of the src array to the dest array.

CMP_ARR_COL

Syntax $bb \leftarrow \text{CMP_ARR_COL} (\text{range } 1, \text{range } 2)$

Preconditions $\text{range } 1$ and $\text{range } 2$ must belong to $1..LAUC_maxobj$.

Output bb is a Boolean element that indicates whether array ranges 1 and 2 are identical.

Comparison between the two arrays.

SET_ARR_COL

Syntax $\text{SET_ARR_COL} (\text{range}, ii, jj, vv)$

Preconditions range belonging to $\text{dom}(\text{arr_col})$, i.e. it corresponds to the index of a previously created array. ii and jj are in $1..LAUC_maxidx$, jj must be different from MAXINT . vv is in $LAUC_minval..LAUC_maxval$.

The value vv is copied to the range array for all indexes between ii and jj . If $ii > jj$, the array remains unchanged.

PCOP_ARR_COL

Syntax $\text{PCOP_ARR_COL} (\text{dest}, \text{idx_dst}, \text{src}, ii, jj)$

Preconditions dest and src are elements that are different from $1..LAUC_maxobj$, corresponding to arrays already created. $ii..jj$ is a non blank interval of $0..LAUC_maxidx$ and $jj \neq \text{MAXINT}$. $\text{idx_dst}.. \text{idx_dst} + jj - ii$ is an interval of $0..LAUC_maxidx$.

The $ii..jj$ part in the src array is copied to the $\text{idx_dst}.. \text{idx_dst} + jj - ii$ part of the dst array.

PCMP_ARR_COL

Syntax $\text{idx}, \text{bb} \leftarrow \text{PCMP_ARR_COL} (\text{nn2}, \text{idx2}, \text{nn1}, \text{ii}, \text{jj})$

Preconditions nn1 and nn2 are elements that are different from 1..LAUC_maxobj and correspond to arrays already created. $\text{ii}..\text{jj}$ is a non blank interval of 0..LAUC_maxidx. $\text{idx2}..\text{idx2} + \text{jj} - \text{ii}$ is an interval of 0..LAUC_maxidx.

Outputs bb is a BOOL. idx is in $\text{ii}..\text{jj}$.

The $\text{ii}..\text{jj}$ part in array nn1 is compared to part $\text{idx2}..\text{idx2} + \text{jj} - \text{ii}$ in array nn2 . bb is FALSE if the two parts are identical, TRUE if not. In this case, idx is the index of the first element that is different from $\text{ii}..\text{jj}$.

IMPORTS REQUIRED

(instances to import as the implementation tree for this library machine sees them with SEES) BASIC_ARITHMETIC, BASIC_BOOL.